

# OCAML

## ☞ nucleo funzionale puro

- funzioni (ricorsive)
- tipi e pattern matching
- primitive utili: liste
- trascriviamo pezzi della semantica operativa

## ☞ componente imperativo

- variabili e assegnamento
- primitive utili: arrays

## ☞ moduli

# Espressioni pure

Objective Caml version 2.00+3/Mac1.0a1

```
# 25;;
```

```
- : int = 25
```

```
# true;;
```

```
- : bool = true
```

```
# 23 * 17;;
```

```
- : int = 391
```

```
# true & false;;
```

```
- : bool = false
```

```
# 23 * true;;
```

This expression has type bool but is here used with type int

```
# if 2 = 3 then 23 * 17 else 15;;
```

```
- : int = 15
```

```
# if 2 = 3 then 23 else true;;
```

This expression has type bool but is here used with type int

# Funzioni

```
# function x -> x + 1;;  
- : int -> int = <fun>  
# (function x -> x + 1) 3;;  
- : int = 4  
# (function x -> x + 1) true;;  
This expression has type bool but is here used with type int  
# function x -> x;;  
- : 'a -> 'a = <fun>  
# function x -> function y -> x y;;  
- : ('a -> 'b) -> 'a -> 'b = <fun>  
# (function x -> x) 2;;  
- : int = 2  
# (function x -> x) (function x -> x + 1);;  
- : int -> int = <fun>  
# function (x, y) -> x + y;;  
- : int * int -> int = <fun>  
# (function (x, y) -> x + y) (2, 33);;  
- : int = 35
```

# Let binding

```
# let x = 3;;  
val x : int = 3  
# x;;  
- : int = 3  
# let y = 5 in x + y;;  
- : int = 8  
# y;;  
Unbound value y  
# let f = function x -> x + 1;;  
val f : int -> int = <fun>  
# f 3;;  
- : int = 4  
# let f x = x + 1;;  
val f : int -> int = <fun>  
# f 3;;  
- : int = 4  
# let fact x = if x = 0 then 1 else x * fact(x - 1) ;;  
Unbound value fact
```

# Let rec binding

```
# let rec fact x = if x = 0 then 1 else x * fact(x - 1) ;;  
val fact : int -> int = <fun>  
# fact (x + 1);;  
- : int = 24
```

# Tipi 1

```
# type ide = string;;  
type ide = string  
# type expr = | Den of ide | Val of ide | Fun of ide * expr  
| Plus of expr * expr | Apply of expr * expr;;  
type expr =  
| Den of ide  
| Val of ide  
| Fun of ide * expr  
| Plus of expr * expr  
| Apply of expr * expr
```

$E ::= I \mid \text{val}(I) \mid \text{lambda}(I, E_1) \mid \text{plus}(E_1, E_2) \mid \text{apply}(E_1, E_2)$

```
# Apply(Fun("x",Plus(Den "x", Den "x")), Val "y");;  
- : expr = Apply (Fun ("x", Plus (Den "x", Den "x")), Val "y")
```

# Tipi 2

```
# type eval = Int of int | Bool of bool | Efun of expr
  | Unbound;;
type eval = | Int of int | Bool of bool | Efun of expr |
  Unbound
# type env = ide -> eval;;
type env = ide -> eval
# let bind (rho, i, d) =
  function x -> if x = i then d else rho x;;
val bind : (ide -> eval) * ide * eval -> (ide -> eval) = <fun>
```

- $env = IDE \rightarrow eval$
- $eval = [ int + bool + fun ]$

# Tipi 3

```
# type com = Assign of ide * expr | Ifthenelse of expr *  
com list * com list | While of expr * com list;;
```

```
type com =
```

```
| Assign of ide * expr  
| Ifthenelse of expr * com list * com list  
| While of expr * com list
```

```
C ::= ifthenelse(E, C1, C2) | while(E, C1) | assign(I, E) | cseq(C1, C2)
```

```
# While(Den "x", [Assign("y", Plus(Val "y", Val "x"))]);;  
- : com = While (Den "x", [Assign ("y", Plus (Val "y", Val  
"x"))])
```



# Un tipo primitivo utile: le liste

```
# let l1 = [1; 2; 1];;  
val l1 : int list = [1; 2; 1]  
# let l2 = 3 :: l1;;  
val l2 : int list = [3; 1; 2; 1]  
# let l3 = l1 @ l2;;  
val l3 : int list = [1; 2; 1; 3; 1; 2; 1]  
# List.hd l3;;  
- : int = 1  
# List.tl l3;;  
- : int list = [2; 1; 3; 1; 2; 1]  
# List.length l3;;  
- : int = 7
```

# Tipi e pattern matching

```
type expr =
  | Den of ide
  | Fun of ide * expr
  | Plus of expr * expr
  | Apply of expr * expr
type eval = | Int of int | Bool of bool | Efun of expr | Unbound
type env = ide -> eval
 $\mathcal{E}(I, \rho) = \rho(I)$ 
 $\mathcal{E}(\text{plus}(E_1, E_2), \rho) = \mathcal{E}(E_1, \rho) + \mathcal{E}(E_2, \rho)$ 
 $\mathcal{E}(\text{lambda}(I, E_1), \rho) = \text{lambda}(I, E_1)$ 
 $\mathcal{E}(\text{apply}(E_1, E_2), \rho) = \text{applyfun}(\mathcal{E}(E_1, \rho), \mathcal{E}(E_2, \rho), \rho)$ 
 $\text{applyfun}(\text{lambda}(I, E_1), d, \rho) = \mathcal{E}(E_1) [\rho / I \leftarrow d]$ 
# let rec sem (e, rho) = match e with
  | Den i -> rho i
  | Plus(e1, e2) -> plus(sem (e1, rho), sem (e2, rho))
  | Fun(i, e) -> Efun(Fun(i, e))
  | Apply(e1, e2) -> match sem(e1, rho) with
    | Efun(Fun(i, e)) -> sem(e, bind(rho, i, sem(e2, rho)))
    | _ -> failwith("wrong application");;
val sem : expr * env -> eval = <fun>
```

# Variabili e frammento imperativo

```
# let x = ref(3);;
val x : int ref = {contents=3}
# !x;;
- : int = 3
# x := 25;;
- : unit = ()
# !x;;
- : int = 25
# x := !x + 2; !x;;
- : int = 27
```

# Un tipo primitivo mutabile: l'array

```
# let a = [| 1; 2; 3 |];;
val a : int array = [|1; 2; 3|]
# let b = Array.make 12 1;;
val b : int array = [|1; 1; 1; 1; 1; 1; 1; 1; 1; 1; 1; 1|]
# Array.length b;;
- : int = 12
# Array.get a 0;;
- : int = 1
# Array.get b 12;;
Uncaught exception: Invalid_argument("Array.get")
# Array.set b 3 131;;
- : unit = ()
# b;;
- : int array = [|1; 1; 1; 131; 1; 1; 1; 1; 1; 1; 1; 1|]
```

# Moduli: interfaccie

```
# module type PILA =  
  sig  
    type 'a stack          (* abstract *)  
    val emptystack : 'a stack  
    val push : 'a stack -> 'a -> 'a stack  
    val pop : 'a stack -> 'a stack  
    val top : 'a stack -> 'a  
  end;;  
module type PILA =  
  sig  
    type 'a stack  
    val emptystack : 'a stack  
    val push : 'a stack -> 'a -> 'a stack  
    val pop : 'a stack -> 'a stack  
    val top : 'a stack -> 'a  
  end
```

# Moduli: implementazione

```
# module SpecPila: PILA =
  struct
    type 'a stack = Empty | Push of 'a stack * 'a
    let emptystack = Empty
    let push p a = Push(p,a)
    let pop p = match p with
      | Push(p1, _) -> p1
    let top p = match p with
      | Push(_, a) -> a
  end;;
module SpecPila : PILA
```

# Il linguaggio didattico

☞ le cose semanticamente importanti di OCAML, meno

- tipi (e pattern matching)
- moduli
- eccezioni