Laurea Magistrale in INFORMATICA Principi di Linguaggi di Programmazione Compiler Techniques

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Regular Ėxam V - July 18th 2014

(Available Time: 2 hours. Mandatory: In each exercise get at least, one half of the assigned points)

Exercise 1 (pts 10). Let $E1 = ((b | a) c^*)^*$ and $E2 = (a | b | c)^*$.

- a) Using dotted automata, provide a proof of the equivalence of E1 and E2;
- b) Apply to E1, the method for transforming DFSA into right-linear grammars.
- c) Show the LL(1) parsing table of the grammar You obtain in (b).

Exercise 2 (pts 10). Let L be the following language:

- $L = \{a^{n1} c \dots a^{nk} c b^m d^{n-m} | k \ge 1, n = n1 + \dots + nk, (\forall i: 1 \le i \le k) ni > 0\}$
- (a) Provide a LR grammar G for L;
- (b) Provide the collection Coll(1) of G
- (c) Looking at the collection, answer:

[1] is \overline{G} an LR(1) grammar?

[2] is G a LALR(1) grammar?

[3] is G a SLR(1) grammar.

Exercise 3 (pts 10). Extend the language "Semplice" with the multiple assignment defined below

Cmd ::= MUpd MUpd ::= Ide IdeList := Exp ExpList IdeList ::= , Ide Idelist | ε

Exp ::= , Exp Explist | ε

Multiple assignment requires that the number of the identifiers on the left hand of := is the same of the number of expression on the right hand. Moreover, it behaves like any arbitrary sequence of assignments "ide_i := \exp_i " where $1 \le i \le n$, ide_i (resp. \exp_i) is the i-th identifier (resp. expression) of the list on the left (resp. right) hand of :=.

- 1) Provide a translation scheme of the productions above, with attributes that:
 - a) check whether or not the requirement on the size of the lists is satisfied;
 - b) produce, by side-effects, the 3AC code of the multiple;
- 2) Apply the translation in 1.b) to the multiple assignment:

x,y := 10,x+y

and show the generated 3AC code.