

Assigning Types to Program Structures of a Strongly Typed Language

Types

Types = Basic-Types + N + Types X Types \rightarrow Types

Rules

$$\frac{\Gamma \vdash i:t \quad \Gamma \vdash e:T}{\Gamma \vdash i:=e:N}$$

$$\frac{\Gamma \vdash e:\text{boolean} \quad \Gamma \vdash Cs:N}{\Gamma \vdash \text{while } e \text{ do } Cs:N}$$

$$\frac{\Gamma \vdash C:N \quad \Gamma \vdash Cs:N}{\Gamma \vdash C;Cs:N}$$

$$\frac{\Gamma_1, i:T, \Gamma_2 \vdash i:T \quad \Gamma \vdash e_1:T_1 \quad \Gamma \vdash e_2:T_2 \quad \Gamma \vdash \text{op}:T_1 \times T_2 \rightarrow T}{\Gamma \vdash \text{op}(e_1, e_2):T}$$

$$\frac{\Gamma, i:T_i, j:T_j \vdash e:T}{\Gamma \vdash \text{fun}(i:T_i, j:T_j)\{e\} : T_i \times T_j \rightarrow T}$$

Attributes:

r:- type of all the program structures but expressions

-synthesized of P, C, Cs, A, W

type:- type of expression

-synthesized of E, F, E', F', T, ide

in:- type expression of the left brother on the left

-inherited of E', F'

It is extending the first grammar in its last variant discussed in [variant](#)

Try to do it yourself + then compare your solution with the one in the next slide.

A type checker for Simple Commands

[0] $P ::= Ds Cs$	$P.r := Cs.r$
[1] $P ::= Cs$	$P.r := Cs.r$
[2] $Ds ::= Var Dts$	
[3] $Dts ::= Dt Dts'$	
[4] $Dts' ::= ; Dt Dts'$	
[5] $Dts' ::= \epsilon$	
[6] $Dt ::= ide O$	$addtype(ide.entry, O.t)$
[7] $O_1 ::= , ide O_2$	$addtype(ide.entry, O_2.t); O_1.t := O_2.t$
[8] $O ::= : Y$	$O.t := Y.t$
[9] $Cs_1 ::= ; C Cs_2$	$Cs_1.r := if (C.r=N \& Cs_1.r=N) then N else \perp$
[10] $Cs ::= \epsilon$	$Cs.r := N$
[11] $C ::= A$	$C.r := A.r$
[12] $C ::= W$	$C.r := W.r$
[13] $A ::= ide := E$	$A.r := if (ide.type=E.type \& ide.type \neq \perp)$ $then N else \perp$
[14] $W ::= while E do C Cs endw$	$W.r := if (E.type=boolean \& C.r=Cs.r)$ $then C.r else \perp$
[24] $Y ::= boolean$	$Y.t := boolean$
[25] $Y ::= integer$	$Y.t := integer$
[26] $Y ::= file$	$Y.t := file$

A type checker for Simple Expressions

[15] $E ::= F E'$	{ $E'.in = F.type$; $E.type := E'.type$;
[16] $E'_1 ::= op-l F E'_2$	{ $match\ t1xt2 \rightarrow t = op-l.type$; if ($E'_1.in = t1 \ \& \ F.type = t2$) then { $E'_2.in = t$; $E'_1.type := E'_2.type$ }; else { $E'_2.in = \perp$; $E'_1.type := \perp$ }; }
[17] $E' ::= \epsilon$	{ $E'.type = E'.in$ }
[18] $F ::= T F'$	{...}
[19] $F' ::= op-h T F'$	{...}
[20] $F' ::= \epsilon$	{...}
[21] $T ::= num$	{ $T.type = integer$;
[22] $T ::= ide$	{ $T.type = ide.type$;
[23] $T ::= (E)$	{ $T.type = E.type$;