Positive matrices

We write $P \ge Q$ if $P_{ij} \ge Q_{ij}$ for all i, j. We write P > Q if $P_{ij} > Q_{ij}$ for all i, j.

(Also, $P \ge 0, P > 0.$)

If $A \ge 0, B \ge 0$, then $AB \ge 0$.

Strict positivity trick

If A > 0, $\mathbf{v} \ge \mathbf{0}$, then $A\mathbf{v} > \mathbf{0}$, unless $\mathbf{v} = \mathbf{0}$. If A > 0, $\mathbf{u} \ge \mathbf{w}$, then $A\mathbf{u} > A\mathbf{w}$, unless $\mathbf{u} = \mathbf{w}$. (Proof: set $\mathbf{v} = \mathbf{u} - \mathbf{w}$ above.)

Perron-Frobenius theorem

Theorem

Let P > 0 be a square matrix. Then,

- 1. *P* has an eigenvalue $\lambda > 0$ with eigenvector $\mathbf{v} > \mathbf{0}$, i.e., $P\mathbf{v} = \lambda \mathbf{v}$.
- 2. λ is the largest eigenvalue in absolute value.
- 3. λ has multiplicity 1, and it is the only eigenvalue with eigenvector $\mathbf{v} > 0$ (up to multiples).

If $P \ge 0, 1., 2., 4$. hold with \ge instead of >, but in many cases also the original statement holds (unless P is associated to a 'disconnected' or 'periodic' graph — we'll see it later.)

 $\lambda=\rho(P), \mathbf{v}$ are called 'the Perron eigenvalue/eigenvector' of P. Also,

4. (monotonicity) If $0 \le P \le Q$, then $0 \le \rho(P) \le \rho(Q)$.

Proof (just a sketch while you are still awake)

We say that *P* stretches a vector $\mathbf{v} \ge \mathbf{0}$ by a factor k > 0 if $P\mathbf{v} \ge k\mathbf{v}$ (and *k* is the largest possible). Example

$$P = \begin{bmatrix} 0.2 & 0.7 & 0.1 \\ 0.3 & 0.4 & 0.5 \\ 0.1 & 0.7 & 0.2 \end{bmatrix}, \quad \mathbf{v} = \begin{bmatrix} 1 \\ 2 \\ 3 \end{bmatrix}, \quad P\mathbf{v} = \begin{bmatrix} 1.9 \\ 2.6 \\ 2.1 \end{bmatrix}$$

P stretches **v** by a factor $\min(\frac{1.9}{1}, \frac{2.6}{2}, \frac{2.1}{3}) = 0.7$. Indeed, 0.7 $\begin{bmatrix} 1\\2\\3 \end{bmatrix} \leq \begin{bmatrix} 1.9\\2.6\\2.1 \end{bmatrix}$.

Proof (cont.)

By the strict positivity trick,

$$k\mathbf{v} \leq P\mathbf{v} \implies kP\mathbf{v} < P(P\mathbf{v}),$$

i.e., *P* stretches $P\mathbf{v}$ by strictly more than it stretches \mathbf{v} ; unless $k\mathbf{v} = P\mathbf{v}$.

Hence if we take the vector \mathbf{v}_{\max} which has the maximum stretch factor k_{\max} , then it must be the case that $k_{\max}\mathbf{v}_{\max} = P\mathbf{v}_{\max}$. This proves 1.

2. can be proved by taking another eigenvector ${\bf w}$ and considering the stretch factor of $|{\bf w}|.$

Proof (cont.)

3. can be proved by contradiction: assume *P* has another eigenpair $P\mathbf{w} = \mu\mathbf{w}$ with $\mathbf{w} > \mathbf{0}$ in addition to the Perron one $P\mathbf{v} = \lambda\mathbf{v}$.

The transposed matrix P^T has the same eigenvalue λ : $P^T \mathbf{u} = \lambda \mathbf{u}$, i.e., $\mathbf{u}^T P = \lambda \mathbf{u}^T$. (This is called sometimes a left eigenvector of P). By Part 1 of the theorem, $\mathbf{u} > \mathbf{0}$.

Compute in two ways:

$$\lambda \mathbf{u}^T \mathbf{w} = (\mathbf{u}^T P) \mathbf{w} = \mathbf{u}^T P \mathbf{w} = \mathbf{u}^T (P \mathbf{w}) = \mu \mathbf{u}^T \mathbf{w}.$$

4. follows by the fact that Q stretches the Perron vector **v** of P by a factor at least λ .

M-matrices

Similarly, matrices with sign pattern

$$\begin{bmatrix} + & - & - & - \\ - & + & - & - \\ - & - & + & - \\ - & - & - & + \end{bmatrix}$$

(possibly with zeros) can be seen as sI - P, for some scalar $s \ge 0$ and matrix $P \ge 0$, and some results on their eigenvalues can be derived from this (just a heads-up).

Markov chains

TL;DR: finite state automaton + transition probabilities.



At every (discrete) 'time step', we follow an arrow exiting from the current state.

Markov property: transition probabilities do not depend on 'what happened earlier / where I am coming from'.

$$\mathbb{P}[s_k = j \mid s_1 = i_1, s_2 = i_2, \dots, s_{k-1} = i_{k-1}] = \mathbb{P}[s_k = j \mid s_{k-1} = i_{k-1}].$$

Homogeneity: transition probabilities do not change with step k.

Markov chains



Transition probability matrix: $P_{ij} = \mathbb{P}[\text{transition } i \rightarrow j]$, e.g.,

$$P = \begin{bmatrix} 0.4 & 0.6\\ 0.7 & 0.3 \end{bmatrix}.$$

P is row-stochastic, i.e., $\sum_{j} P_{ij} = 1$ for each row *i*. Or, in other words, $P\mathbf{1} = \mathbf{1}$ for the vector $\mathbf{1}$ of all ones (Perron vector with eigenvalue 1!)

Markov chain and linear algebra

Key idea: computing transition probabilities = matrix multiplication.

If $\pi = [\pi_1, \pi_2, \dots, \pi_n]$ (often with $\pi \cdot \mathbf{1} = 1$) contains probabilities of being in each state at a certain time t, then πP contains the probabilities at time t + 1.

Proof:

$$\mathbb{P}[s_{t+1}=j]=\sum_{i}\mathbb{P}[s_t=i]P_{ij}=\pi P.$$

Markov chain and matrix products

Starting from a certain initial probability π , the probability of observing a transition with probability P_1 first, then one with probability P_2, \ldots , then one with probability P_k , is

$$\pi P_1 P_2 \dots P_k.$$

This is true not only when π, P_1, \ldots, P_k are scalar, but also when they are vectors / matrices.

Hitting probabilities

Example: hitting probabilities.

Markov chain with a set of n_B 'bad' states, and n_G 'good' states. Start out in a good state (with probabilities

resp. $\pi = [\pi_1, p_2, \dots, p_{n_G}]$; what is the probability of reaching ('hitting') a bad state? Which one of them is reached first?

$$P = \begin{bmatrix} P_{GG} & P_{GB} \\ P_{BG} & P_{BB} \end{bmatrix}$$

To reach a bad state, we can either:

- transition directly to a bad state, with probability πP_{GB} ;
- transition to a good state once, then to a bad state, πP_{GG}P_{GB};
- transition through 2 good states, then to a bad state, $\pi P_{GG}^2 P_{GB}$; etc.

Total probability:

$$\pi(I + P_{GG} + P_{GG}^2 + P_{GG}^3 + \dots)P_{GB} = \pi(I - P_{GG})^{-1}P_{GB}.$$

Hitting probabilities

$$F_{GB} = (I + P_{GG} + P_{GG}^2 + P_{GG}^3 + \dots)P_{GB} = (I - P_{GG})^{-1}P_{GB}.$$

 $(F_{GB})_{ij}$ gives the probabilities that the Markov chain first enters the set of bad states *B* in its *j*th state, starting from the *i*th good state.

Remarks:

- ► The formula $(I + M + M^2 + M^3 + ...) = (I M)^{-1}$ holds for each square matrix M with $\rho(M) < 1$.
- ▶ P_{GG} satisfies $P_{GG}\mathbf{1} \leq \mathbf{1}$, and the \leq is not an equal (unless $P_{GB} = 0$). So by monotonicity $\rho(P_{GG}) < 1$.

Also, one can prove that the 'mean hitting time' is $(I - P_{GG})^{-1}\mathbf{1}$.

Example: a game with coins

We toss a coin repeatedly. I win if three consecutive tosses give THH, you win if you get HHT first. Who is at an advantage?

Set up transition matrix over all 8 possible sequences of 3 tosses, compute hitting probabilities for the set of bad states THH, HHT.

From the initial set $\pmb{\pi} = [1/8, 1/8, \dots, 1.8]$, we get

 $\pi_B + \pi_G F_{GB} = \dots$

Censoring

Censoring: rewrite transition history 'pretending the set G does not exist'.

$$b_1, g_1, b_2, b_2, g_2, g_3, g_1, g_2, b_1, b_3, g_1, g_2, b_2, \dots$$

becomes

 $b_1, b_2, b_2, b_1, b_3, b_2, \ldots$

Transition matrix of the censored chain $n_B \times n_B$:

$$P_{BB} + P_{BG}(I - P_{GG})^{-1}P_{GB}$$

Interpretation: we either transition from B to B directly, or we transition to G, stay inside it for 0, 1, 2, ... time steps, and then get back out to B.

Interesting interpretation: Gaussian elimination on $I - P \iff$ censoring states 1, 2, 3, . . . in sequence.

Stationary probabilities

Suppose P > 0 for now (we'll see what changes if there are zero entries).

Theorem

For any initial probabilities π , the probabilities $\pi P, \pi P^2, \ldots, \pi P^k, \ldots$ of being in each state after k steps converge to a fixed vector μ when $k \to \infty$. This vector is the left eigenvector with eigenvalue $\rho(P) = 1$, i.e., the Perron vector of P^T .

Proof 1: this is simply the power method on P^{T} .

Proof 2: if $P = V\Lambda V^{-1}$ is diagonalizable, recall that all other eigenvalues apart from $1 = \lambda_1 = \rho(P)$ have modulus $|\lambda_i| < 1$, hence (cont.)

$$\pi (V\Lambda V^{-1})^k = \pi V egin{bmatrix} \lambda_1^k & & & \ & \lambda_2^k & & \ & & \ddots & \ & & & \ddots & \ & & & & \lambda_n^k \end{bmatrix} V^{-1}$$

which converges when $k \to \infty$ to

$$\pi (V \wedge V^{-1})^k = \pi V \begin{bmatrix} 1 & & & \\ & 0 & & \\ & & \ddots & \\ & & & 0 \end{bmatrix} V^{-1}$$

which is a multiple of the first row of V^{-1} .

One can prove directly that the first row of V^{-1} is a left eigenvector, exactly like one proves that the columns of V are the right eigenvectors. Otherwise, take limits in both sides of $(\pi P^k)P = \pi P^{k+1}$.

What happens if there are zeros

Some states (called transient states) are visited only a few times and then 'abandoned forever', e.g.,



Key to understand it: doubly connected components of the graph with adjacency matrix *P*.

If a DCC has outgoing edges, eventually they will be taken with probability $1 \implies$ the DCC will be abandoned with probability 1.

Each DCC without outgoing nodes is a possible final class.

Periodic chains

Another edge case: periodic chains. Suppose P is irreducible, but all closed paths have lengths that are multiple of a certain integer d > 1; e.g.,



In general, their transition matrices can be written as



(These may be blocks.)

General Perron-Frobenius

Periodic chains do not have a limit distribution (easy to see also in the example; starting from state 1 one 'loops indefinitely').

Another characterization: a chain is aperiodic if $P^k > 0$ for some k.

Periodic chains have d eigenvalues with modulus 1 (at the dth roots of 1). In particular, 'Perron-Frobenius with strict inequalities' does not hold for them.

The missing hypothesis

The Perron–Frobenius theorem 'with strict inequalities' holds for matrices $P \ge 0$ that have a doubly-connected graph (irreducible chains/matrices) and are aperiodic.

References

Meyer, *Matrix Analysis and Applied Linear Algebra*, chapter 8. — gentler introduction

Berman, Plemmons, *Nonnegative matrices in the mathematical sciences* — more technical; includes a whopping "50 equivalent conditions for a matrix to be a nonsingular M-matrix" (!)