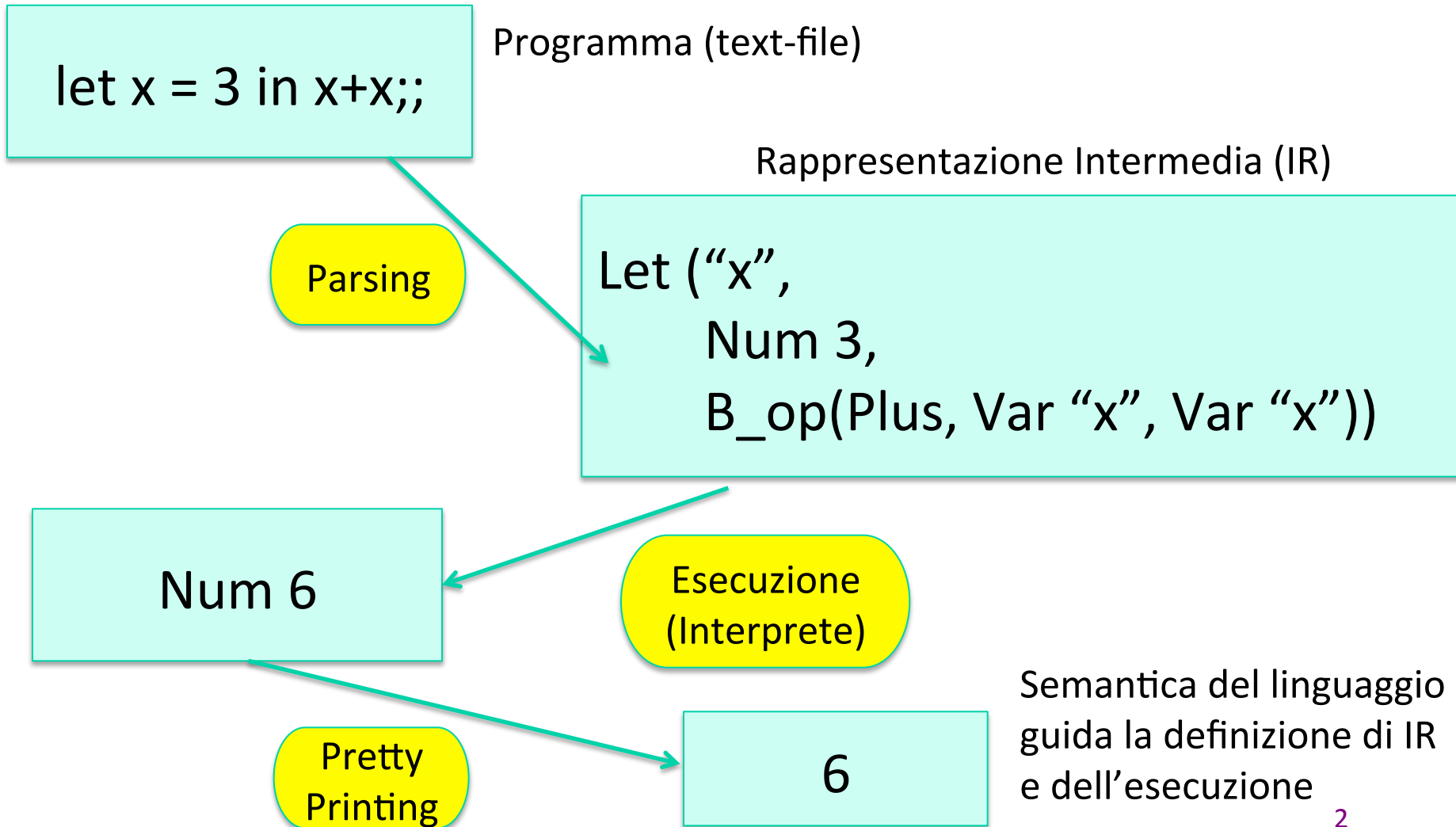

PROGRAMMAZIONE 2

16. Realizzare un interprete in OCaml

La struttura



La struttura nel dettaglio

OCaML Type per descrivere la rappresentazione intermedia

```
type variable = string

type operand = Plus | Minus | Times | ...

type exp =
  Int_e of int
  | Op_e of exp * op * exp
  | Var_e of variable
  | Let_e of variable * exp * exp
```

La struttura nel dettaglio

```
type variable = string
```

```
type operand = Plus | Minus | Times | ...
```

```
type exp =
```

```
  | Int_e of int
```

```
  | Op_e of exp * op * exp
```

```
  | Var_e of variable
```

```
  | Let_e of variable * exp * exp
```

***Rappresentazione di
"3 + 17"***

```
let e1 = Int_e 3
```

```
let e2 = Int_e 17
```

```
let e3 = Op_e (e1, Plus, e2)
```

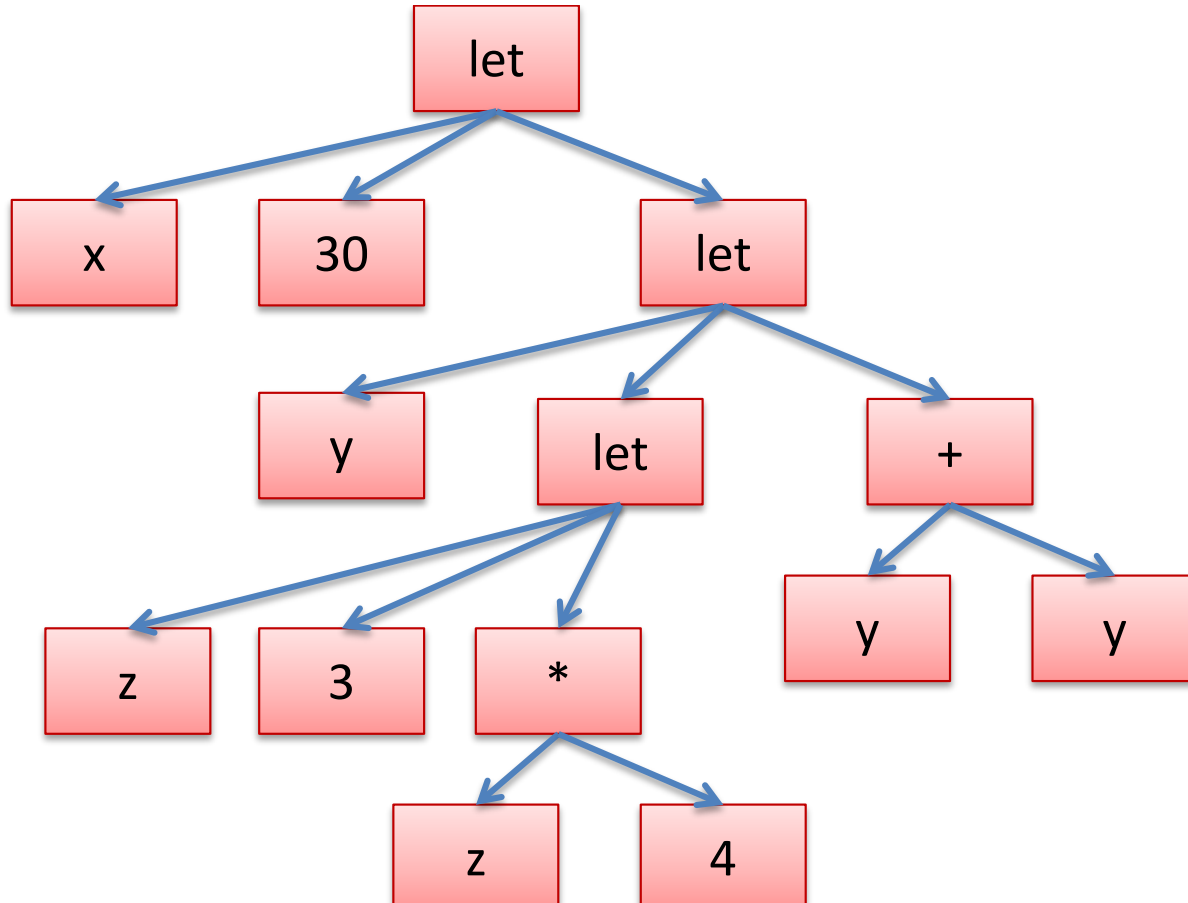
```
let x = 30 in
let y =
  (let z = 3 in z*4) in
  y+y;;
```

Programma OCaml

Exp

```
Let_e("x", Int_e 30,
  Let_e("y",
    Let_e("z", Int_e 3, Op_e(Var_e "z", Times, Int_e 4)),
    Op_e(Var_e "y", Plus, Var_e "y"))
```

AST



Variabili: dichiarazione e uso

```
type variable = string
```

```
type exp =
```

```
    Int_e of int
```

```
    | Op_e of exp * op * exp
```

```
    | Var_e of variable
```

```
    | Let_e of variable * exp * exp
```

Runtime: operazione di supporto

eval_op : exp -> op -> exp -> exp

substitute : exp -> variable -> exp -> exp

Variabili: dichiarazione e uso

```
type variable = string
```

```
type exp =
```

```
  Int_e of int
```

```
  | Op_e of exp * op * exp
```

```
  | Var_e of variable
```

```
  | Let_e of variable * exp * exp
```

**Uso di
una
variabile**

Variabili: dichiarazione e uso

```
type variable = string
```

```
type exp =
```

```
  Int_e of int
```

```
  | Op_e of exp * op * exp
```

```
  | Var_e of variable
```

```
  | Let_e of variable * exp * exp
```

**Uso di
una
variabile**

**Dichiarazione
di variable**

L'interprete

RTS

eval_op : exp -> op -> exp -> exp
substitute : exp -> variable -> exp -> exp

```
let rec eval (e : exp) : exp =  
  match e with  
  | Int_e _ ->  
  | Op_e(e1,op,e2) ->  
  | Let_e(x,e1,e2) ->
```

L'interprete

RTS

eval_op : exp -> op -> exp -> exp

substitute : exp -> variable -> exp -> exp

let rec eval (e : exp) : exp =
 match e with

Int_e _ -> **e** (* Int_e i -> Int_e i *)
 | Op_e(e1,op,e2) ->
 | Let_e(x,e1,e2) ->

L'interprete

RTS

eval_op : exp -> op -> exp -> exp
substitute : exp -> variable -> exp -> exp

let rec eval (e : exp) : exp =
 match e with
 | Int_e _ -> e
 | Op_e(e1,op,e2) -> **let v1 = eval e1 in**
 let v2 = eval e2 in
 eval_op v1 op v2
 | Let_e(x,e1,e2) ->

L'interprete

RTS

eval_op : exp -> op -> exp -> exp

substitute : exp -> variable -> exp -> exp

let rec eval (e : exp) : exp =
 match e with

Int_e _ -> e

| Op_e(e1,op,e2) -> **let v1 = eval e1 in**
 let v2 = eval e2 in
 eval_op v1 op v2

| Let_e(x,e1,e2) -> **let v1 = eval e1 in**
 let e2' = substitute v1 x e2 in
 eval e2'

L'interprete

RTS

eval_op : exp -> op -> exp -> exp

substitute : exp -> variable -> exp -> exp

let rec eval (e : exp) : exp =

match e with

Int_e _ -> e

| Op_e(e1,op,e2) -> **eval_op eval e1 op eval e2**

| Let_e(x,e1,e2) -> **let v1 = eval e1 in**

**let e2' = substitute v1 x e2 in
eval e2'**

L'interprete

RTS

eval_op : exp -> op -> exp -> exp

substitute : exp -> variable -> exp -> exp

let rec eval (e : exp) : exp =
 match e with

 Int_e _ -> e

| Op_e(e1,op,e2) -> **eval_op** eval e1 **op** eval e2

| Let_e(x,e1,e2) -> **let v1 = eval e1 in**

let e2' = substitute v1 x e2 in
eval e2'

Come avviene la
 valutazione?

Si usa **let** per
 definirne l'ordine

L'interprete

RTS

eval_op : exp -> op -> exp -> exp

substitute : exp -> variable -> exp -> exp

**let rec eval (e : exp) : exp =
match e with**

Int_e _ -> e

| Op_e(e1,op,e2) -> eval_op eval e1 op eval e2

| Let_e(x,e1,e2) -> let v1 = eval e1 in

**let e2' = substitute v1 x e2 in
eval e2'**

| Var_e _ -> ???

Non dovremmo incontrare una
variabile – avremmo già dovuta
sostituirla con un valore!!

Questo è un **errore di tipo**

L'interprete

RTS

eval_op : exp -> op -> exp -> exp

substitute : exp -> variable -> exp -> exp

**let rec eval (e : exp) : exp =
match e with**

Int_e _ -> e

| Op_e(e1,op,e2) -> eval_op eval e1 op eval e2

| Let_e(x,e1,e2) -> let v1 = eval e1 in

**let e2' = substitute v1 x e2 in
eval e2'**

| Var_e _ -> Null

Questo complica l'interprete:
matching sui risultati delle
chiamate ricorsive di eval

L'interprete

RTS

eval_op : exp -> op -> exp -> exp

substitute : exp -> variable -> exp -> exp

let rec eval (e : exp) : exp =
match e with

Int_e _ -> e

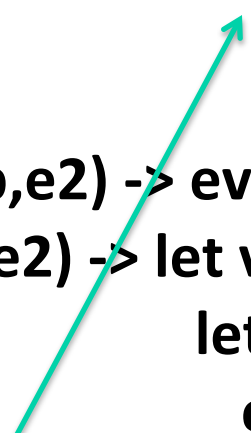
| Op_e(e1,op,e2) -> eval_op eval e1 op eval e2

| Let_e(x,e1,e2) -> let v1 = eval e1 in

let e2' = substitute v1 x e2 in
eval e2'

| Var_e _ -> raise (UnboundVariable x)

Tali eccezioni fanno
 parte del RTS



RTS: eval_op

```
let eval_op (v1:exp) (op:operand) (v2:exp) : exp =  
  match v1, op, v2 with  
    | Int_e i, Plus, Int_e j -> Int_e (i + j)  
    | Int_e i, Minus, Int_e j -> Int_e (i - j)  
    | Int_e i, Times, Int_e j -> Int_e (i * j)  
    | _, _, _ -> raise (BadOp (v1,op,v2))
```

RTS: substitution

```
let substitute (v:exp) (x:variable) (e:exp) : exp =  
  let rec subst (e:exp) : exp =  
    match e with  
    | Int_e _ ->  
    | Op_e(e1,op,e2) ->  
    | Var_e y -> ... use x ...  
    | Let_e (y,e1,e2) -> ... use x ...  
  in subst e
```

RTS: substitution

```
let substitute (v:exp) (x:variable) (e:exp) : exp =  
  let rec subst (e:exp) : exp =  
    match e with  
    | Int_e _ -> e  
    | Op_e(e1,op,e2) ->  
    | Var_e y -> ... use x ...  
    | Let_e (y,e1,e2) -> ... use x ...  
  in subst e
```

RTS: substitution

```
let substitute (v:exp) (x:variable) (e:exp) : exp =  
  let rec subst (e:exp) : exp =  
    match e with  
    | Int_e _ -> e  
    | Op_e(e1,op,e2) -> Op_e(subst e1,op,subst e2)  
    | Var_e y -> ... use x ...  
    | Let_e (y,e1,e2) -> ... use x ...  
  in subst e
```



x, v impliciti

RTS: substitution

```
let substitute (v:exp) (x:variable) (e:exp) : exp =  
  let rec subst (e:exp) : exp =  
    match e with  
    | Int_e _ -> e  
    | Op_e(e1,op,e2) -> Op_e(subst e1,op,subst e2)  
    | Var_e y -> if x = y then v else e  
    | Let_e (y,e1,e2) -> ... use x ...  
  in subst e
```


RTS: substitution

```
let substitute (v:exp) (x:variable) (e:exp) : exp =  
  let rec subst (e:exp) : exp =  
    match e with  
    | Int_e _ -> e  
    | Op_e(e1,op,e2) -> Op_e(subst e1,op,subst e2)  
    | Var_e y -> if x = y then v else e  
    | Let_e (y,e1,e2) -> Let_e (y,subst e1,subst e2)  
  in subst e
```

RTS: substitution

```
let substitute (v:exp) (x:variable) (e:exp) : exp =  
  let rec subst (e:exp) : exp =  
    match e with  
    | Int_e _ -> e  
    | Op_e(e1,op,e2) -> Op_e(subst e1,op,subst e2)  
    | Var_e y -> if x = y then v else e  
    | Let_e (y,e1,e2) -> Let_e (y,subst e1,subst e2)  
  in subst e
```



errore se x=y

RTS: substitution

```

let substitute (v:exp) (x:variable) (e:exp) : exp =
  let rec subst (e:exp) : exp =
    match e with
    | Int_e _ -> e
    | Op_e(e1,op,e2) -> Op_e(subst e1,op,subst e2)
    | Var_e y -> if x = y then v else e
    | Let_e (y,e1,e2) -> Let_e (y,subst e1,
      if x = y then e2
      else subst e2)
  in subst e
  
```

shadow scope

Funzioni



Sintassi

type exp = Int_e of int | Op_e of exp * op * exp
| Var_e of variable | Let_e of variable * exp * exp
| Fun_e of variable * exp | FunCall_e of exp * exp

Sintassi

```
type exp = Int_e of int | Op_e of exp * op * exp
          | Var_e of variable | Let_e of variable * exp * exp
          | Fun_e of variable * exp | FunCall_e of exp * exp
```

La sintassi OCaml **fun x -> e** viene rappresentata come **Fun_e(x, e)**

La chiamata **fact 3** viene rappresentata come **FunCall_e (Var_e “fact”, Int_e 3)**

Esempio (e va cambiato il RTS)

```
let f = fun x -> x + 1 in f 3
```

```
Let_e ("f",  
      Fun_e ("x", Op_e (Var_e "x", Plus, Int_e 1)),  
      FunCall (Var_e "f", Int_e 3)  
      )
```

L'interprete+

RTS

eval_op : exp -> op -> exp -> exp

substitute : exp -> variable -> exp -> exp

let rec eval (e : exp) : exp =
 match e with

:

| **Var_e _** -> raise (UnboundVariable x)

| **Fun_e _** -> e

| **FunCall_e (e1,e2)** ->

match eval e1, eval e2 with

Fun_e (x,e3), v2 -> eval (substitute v2 x e3)

| **_** -> raise (TypeError)

Ricorsione

```
type exp = Int_e of int | Op_e of exp * op * exp
          | Var_e of variable | Let_e of variable * exp * exp
          | Fun_e of variable * exp | FunCall_e of exp * exp
          | Letrec_e of variable * exp * exp
```

```
let rec f = fun x -> f (x + 1) in f 3
```

```
Letrec_e ("f",
  Fun_e ("x",
    FunCall_e (Var_e "f", Op_e (Var_e "x", Plus, Int_e 1))
  ),
  FunCall (Var_e "f", Int_e 3)
)
```

L'interprete++

RTS

eval_op : exp -> op -> exp -> exp
substitute : exp -> variable -> exp -> exp

let rec eval (e : exp) : exp =
 match e with
 :
 | Letrec_e (x,e1,e2) ->
 let e1_unwind =
 substitute (Letrec_e (x,e1,Var_e x)) x e1 in
 eval (Let_e (x,e1_unwind,e2))

Cosa abbiamo imparato?

- OCaML può essere usato come linguaggio per la simulazione della semantica operativa di un linguaggio (incluso se stesso!)
- Vantaggio: simulazione dell'implementazione
- Svantaggio: complicato per le operazioni da effettuare con i tipi di OCaML
 - $Op_e(e1, Plus, e2)$ rispetto a “ $e1 + e2$ ”