

Ambiente

Run-time & Run-time Simulation

Semantica eseguibile

$$\begin{array}{c} env \triangleright e1 \Rightarrow v1 \quad v1 = Funval(Fun(x,e), env1) \\ \frac{env \triangleright e2 \Rightarrow v2 \quad env1[v2 / x] \triangleright e \Rightarrow v}{env \triangleright Apply(e1, e2) \Rightarrow v} \end{array}$$

```
let rec sem ((e: exp), (r: eval env)) =
  match e with
  | ...
  | Fun(x, a) -> Funval(e, r)
  | Apply(e1, e2) -> match sem(e1, r) with
    | Funval(Fun(x, a), r1) ->
      sem(a, bind(r1, x, sem(e2, r)))
    | _ -> failwith("no funct in apply")
```

ricorsione

Rec(i, e) -> makefunrec(i, e, r)

makefunrec(f, Fun(args, body), (r: eval env)) =

let **functional**(rr: eval env) =

bind(r, f, Funval(Fun(args, body), rr)) in

let rec (rfix: string -> eval) =

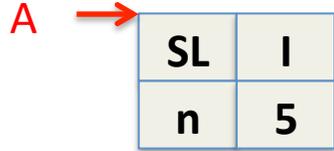
function x -> (functional rfix) x in

Funval(Fun(args, body), rfix)

Un esempio

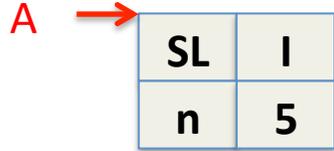
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let n = 5;;  
let h = fun x -> n + x ;;  
let rec f g n = if n = 1 then g(n) else n * f g (n-1);;  
f h 2;;
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Run-time Stack



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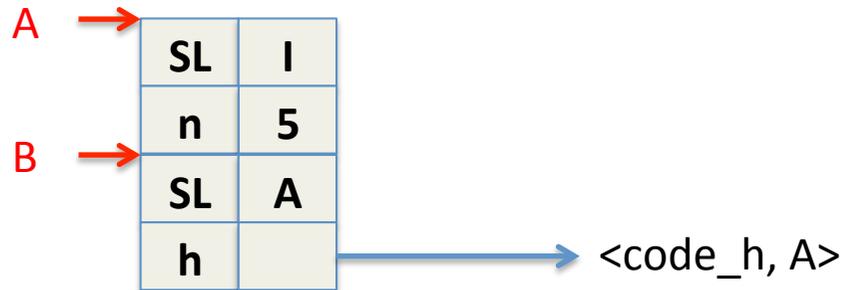
Run-time Stack: simulation



Env_A(n) = 5
Env_A(m) = unbond
for all m ≠ n

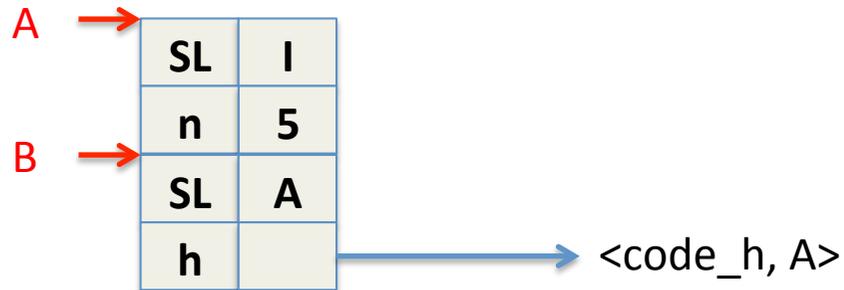
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Run-time Stack



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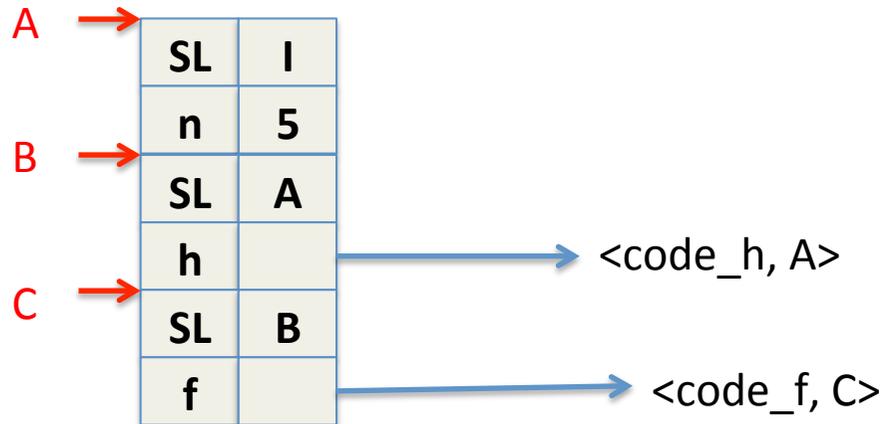
Run-time Stack: simulation



Env_A(n) = 5
Env_A(m) = unbond
for all m != n
Env_B (n) = 5
Env_B(h) = <code_h, Env_A>

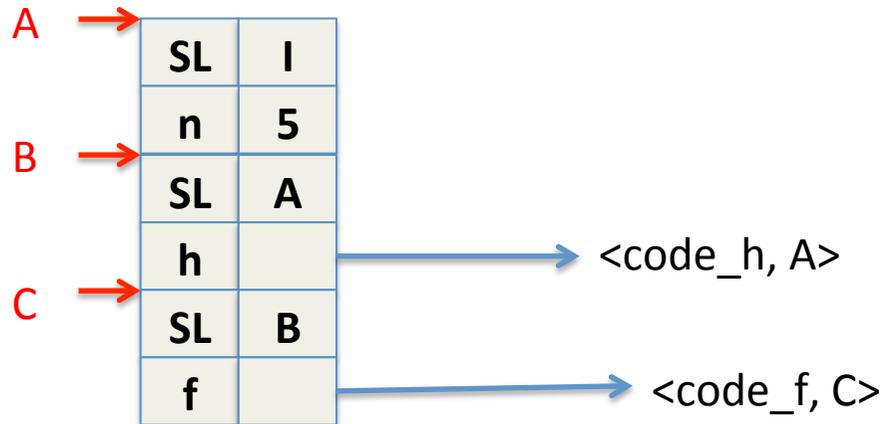
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Run-time Stack: simulation

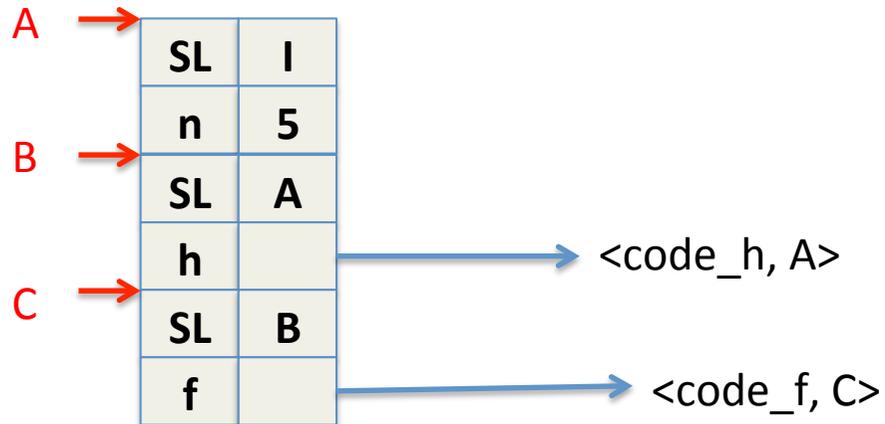


$Env_A(n) = 5$
 $Env_A(m) = \text{unbond}$
 for all $m \neq n$
 $Env_B(n) = 5$
 $Env_B(h) = \langle \text{code}_h, Env_A \rangle$
 $Env_C(f) = \langle \text{code}_f, Env_C \rangle$
 $Env_C(h) = \langle \text{code}_h, Env_A \rangle$
 $Env_C(n) = 5$

```

let n = 5;;
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f h 2;;
    
```

Run-time Stack: simulation



$Env_A(n) = 5$
 $Env_A(m) = \text{unbond}$
 for all $m \neq n$
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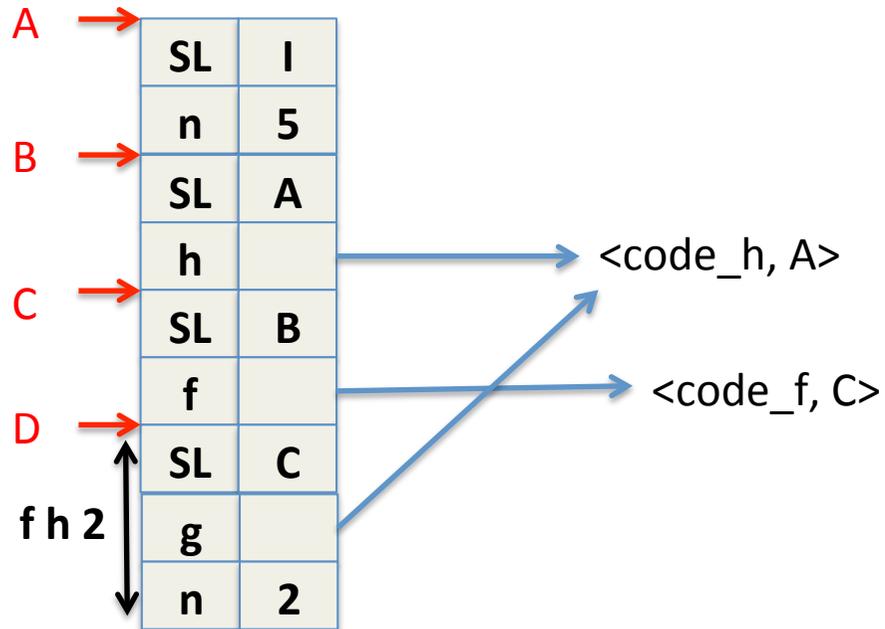
$Env_C(f) = \langle \text{code}_f, Env_C \rangle$
 $Env_C(h) = \langle \text{code}_h, Env_A \rangle$
 $Env_C(n) = 5$

**Definizione ricorsiva:
 makefunrec!!!!**

```

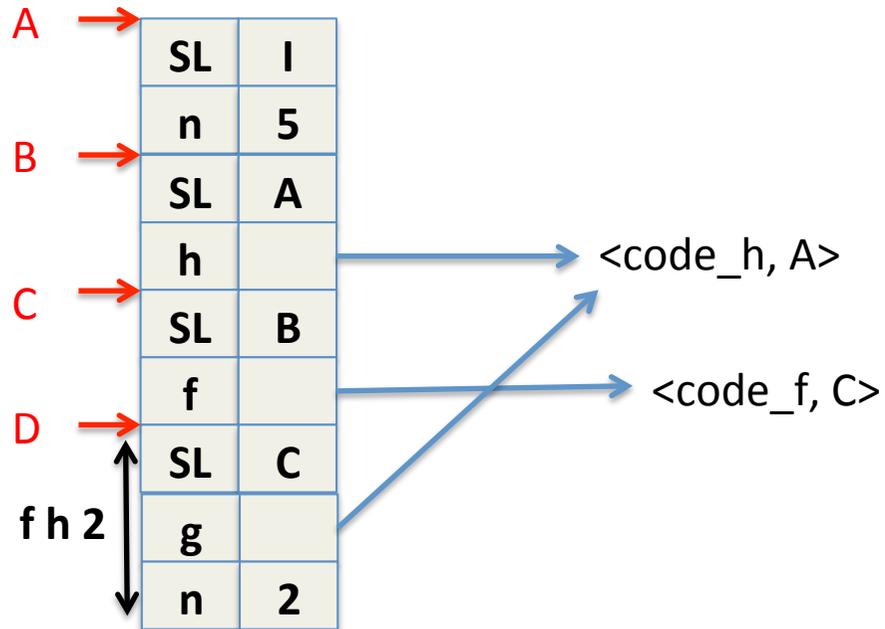
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Run-time Stack



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Run-time Stack: simulation

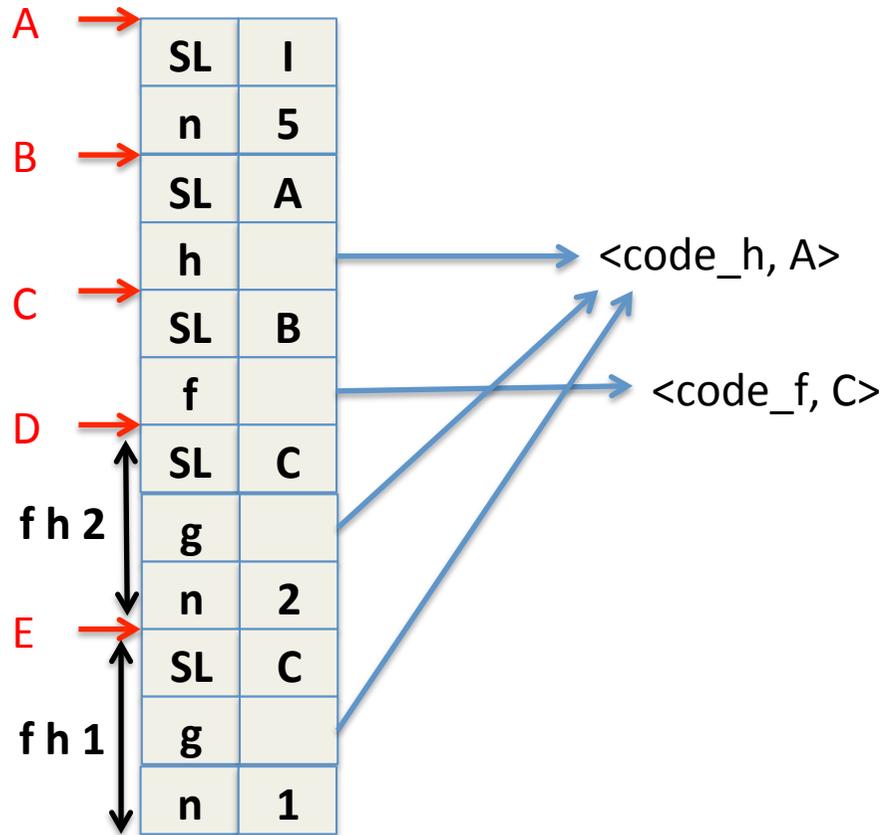


Env_A(n) = 5
 Env_A(m) = unbond
 for all m != n
 Env_B (n) = 5
 Env_B(h) = <code_h, Env_A>
 Env_C(f) = <code_f, Env_C>
 Env_C(h) <code_h, Env_A>
 Env_C(n) = 5
 Env_D(g) = <code_h, Env_A>
 Env_D(n) = 2
 Env_D(f) = <code_f, Env_C>
 Env_D(h) <code_h, Env_A>

```

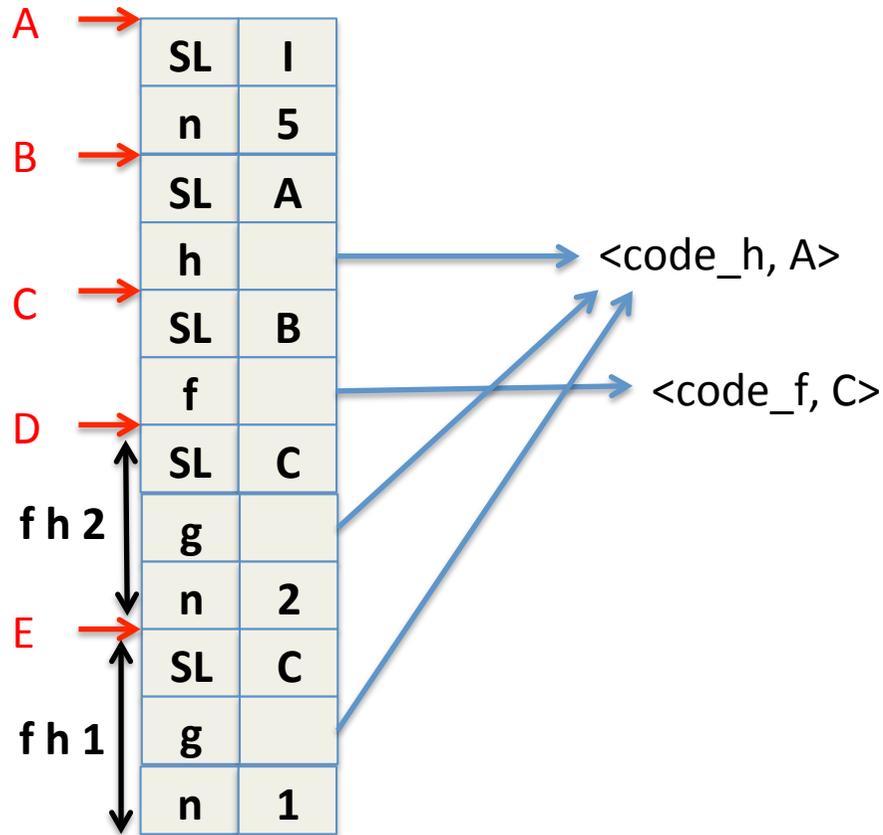
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Run-time Stack



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Run-time Stack: simulation



Env_A(n) = 5
 Env_A(m) = unbond
 for all m != n
 Env_B (n) = 5
 Env_B(h) = <code_h, Env_A>
 Env_C(f) = <code_f, Env_C>
 Env_C(h) <code_h, Env_A>
 Env_C(n) = 5

 Env_D(g) = <code_h, Env_A>
 Env_D(n) = 2
 Env_D(f) = <code_f, Env_C>
 Env_D(h) <code_h, Env_A>

 Env_E(g) = <code_h, Env_A>
 Env_E(n) = 1
 Env_E(f) = <code_f, Env_C>
 Env_E(h) <code_h, Env_A>

```

let n = 5;;
let h = fun x -> n + x ;;
let rec f g n = if n = 1 then g(n) else n * f g (n-1);;
f h 2;;
  
```

Ricorsione

Una simulazione alternativa

Definizione ricorsiva

$$env \triangleright Rec(i, e) \Rightarrow Funval(Rec(i, e), env)$$

$$Rec(i, e) \rightarrow Funval(Rec(i, e), r)$$

REGOLA APPLICAZIONE

$$env \triangleright e1 \Rightarrow v1 \quad v1 = \langle Rec(f, Fun(x, e)), env' \rangle$$
$$env \triangleright e2 \Rightarrow v2$$
$$env[v2 / x][v1 / f] \triangleright e \Rightarrow v$$

$$env \triangleright Apply(e1, e2) \Rightarrow v$$

Ocaml

```
App(e1, e2) -> let e1' = sem e1 env
in (match e1' with
  | Funval(Fun(x, e), env') -> sem e (bindx (eval e2 env) env')
  | Funval(Rec(f, Fun(x,e)), env') ->
      sem e (bind f e1' (bind x (eval e2 env) env'))
  | _ -> raise (TypeError))
```