

Methods as Parameters: A Preprocessing Approach in Java

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Abstract. The paper investigates the use of preprocessing in adding higher order functionalities to Java, that is in passing methods to other methods. The approach is based on a mechanism which offers a restricted, disciplined, form of abstraction that is suitable to the integration of high order and object oriented programming. We discuss how the expressive power of the language is improved. A new syntax is introduced for formal and actual parameters, hence the paper defines a translation that, at preprocessing time, maps programs of the extended language into programs of ordinary Java.

1 Introduction

Higher order programming, *HO*, is considered the main programming methodology of functional languages [Bac78]. In this class of languages, in fact, programs are functions and functions are first class values of the language. This means that functions can be passed as parameters to other functions, returned as result of the computation and furthermore, functions can be values in data structures (i.e. we can have lists, records and arrays of functions). The benefits obtained by HO programming are in the expressiveness of the code, which becomes more concise, clear and well structured and can be reused more easily. These topics are extensively discussed in the literature on functional programming. References, that are a starting point, are [Hud89] and [Wad92].

Though functional languages have never become an effective alternative to the imperative ones, instead limited HO programming features have been added to imperative, first order languages to improve their expressiveness. Examples in this direction are Pascal [JW75], C [KR88] and C++ [Sch95]. Such languages have been defined providing features to allow to pass programs (i.e. procedures or functions) as parameters to other programs. In fact, in this way an abstraction mechanism is added to the language, in particular, programs are generalized with respect to the programs they invoke in the body.

Object-oriented languages improve code reusability and, in a sense, also add higher order features to imperative languages. In fact, in this case, objects are first class values of the language [AC96] and [Cas97]. In Java, objects contain values (instance variables) and methods (instance methods) and can be seen as records [AC96] and interpreted as defining environments, ($Ide \mapsto \{Val \cup Methods\}$), binding selectors to either values or methods, namely *Ide* is the set of selectors that specify object identifiers, *Val* is the set of object values, and *Methods* is the set of object methods. Methods, in this case, are not themselves values, but are contained in objects which are values, hence one method can be *implicitly* passed as a parameter, returned as result and stored in data structures as far as an object containing such a method is passed to, returned from and stored in. In this way, the language provides a kind of higher order, [OW97]. This viewpoint is not

of great help: It leads to programs weighted down from an indirect and tricky use of objects to refer to methods [BO04], and it leads to complicated codes in the attempt to rephrase, in Java, higher order programs already written, for instance, in a functional language.

In [BO05] and [BO06] several approaches to higher order programming in Java are considered. They are characterized by:

- *introspection*, using meta-methods through the package `java.lang.reflect` for meta programming in Java [BO04] and [Bri05];
- *emulation* of a calculus of functions through anonymous inner classes of Java [Set03] or function pointers [McC01];
- *extension* of the language, introducing special entities for function abstractions as in Pizza [OW97] or delegates as in J++ [Cor04].

All the approaches above are valid techniques to support the methodology but none of them seems definitely better than the others. Furthermore, all of them either provide only an indirect way to support the HO methodology or make a neat separation between functions and methods which limits the program expressiveness and makes the use of HO programming a bit tricky. A different solution within the extension approach is discussed in [BO05] and [BO06] where the problem of extending the OO paradigm with method abstraction and method extraction is investigated, in the first order ζ -calculus [AC96], and a way to overcome the conflict originated by combining extraction and subsumption is discussed.

In this paper we discuss an approach to extend Java with mechanisms to pass methods (not arbitrary functions) as parameters and to generalize method invocation. We also discuss how its implementation can be treated at preprocessing time. The main features of the approach are:

- It adds a new kind of parameter: the *parameter method*, `p_method` for short. A `p_method` can be interpreted as defining a function that maps objects (or classes) into instance (respectively, class) methods.
- Since they are functions, `p_methods` could be well defined by resorting to function abstraction. Our approach uses a mechanism much less general which offers a restricted, disciplined, form of function abstraction but suitable to the integration of high order and object oriented programming.
- It enhances code reusability by supplying the language with a higher order based, programming expressiveness.
- All the extensions that it introduces can be implemented by a preprocessing technique that maps programs of the extended language into programs of the ordinary language.

In order to show the benefits of the extensions, in Section 2, we compare the code, in Java, of a program that uses higher order generalization with the code in the extended language. In Section 3 we introduce the language extensions. In Section 4 we formally define the translation \mathcal{E} that maps programs of the extended language into programs of standard Java. The last section concludes the paper.

2 Code reusability: generalization vs. inheritance

We illustrate the reasons and the benefits for extending the language with features for higher order generalization and how these features are integrated with those for object oriented programming, mainly code reusability and inheritance.

2.1 Example of the code development for a class of geometric shapes

Let us consider the development, in Java, of the code for the computation of the lists of the areas and of the perimeters of geometric shapes, equipped with methods to compute the area and the perimeter of the shape [Hud00] [Blo01]. In Figure 1 we start giving a very concise definition for generic *lists* of objects: class `FList` provides `Insert`, `Val` and `Tail` operators. Then we define the class `Shapes` and several subclasses: one for each specific geometric shape, with two obvious methods `Area` and `Perimeter`. The forth class `FListShape` is an extension of `FList`, equipped with two additional methods `FListArea` and `FListPerimeter`, to compute the list of areas and perimeters of a given list of shapes.

2.2 Higher order generalization

As it is clear, examining the code, the method `FListArea` and the method `FListPerimeter` are constituted by the same code except for the name of the method in the invocation. In effect, once we have defined one of the two methods, we would like to obtain the other one, using the mechanisms that the language furnishes for code reusability. Unfortunately in Java these mechanisms are based on class hierarchy and inheritance and they are unable to support this kind of code development. Higher order generalization would provide an adequate support to that form of code development and reusability. In fact it allows both to *generalize* invocation `((Shape)Val()).Area()` into `((Shape)Val()).F()`, where `F` is a functional variable, and to bind `F` to the right method. In the example of Figure 1, this would allow to define a higher order method, `Map`, with a parameter `F`. At each invocation of `Map`, the parameter `F` should be bound to an object operation, namely `Area` or `Perimeter`, which is to be used for `F`. The method `Map` would apply the operation bound to `F` to each object `Val().F()` of the list, and return the list of the computed objects. The invocation of `L.Map(Area)` would compute like `L.FListArea()` while the invocation of `L.Map(Perimeter)` would compute like `L.FListPerimeter()`. Hence, the code for `Map` would be reused for `FListArea()`, `FListPerimeter()` and for any other method that is obtained by instantiating the generalization introduced in the definition of `Map`.

2.3 P_method: A restricted form of function abstraction

We extend Java with a mechanism for higher order generalization of the sort described above. However we cannot simply pass methods as parameters because of the late binding semantics of method overriding [GJSB05]. Instead, we pass a *p_method*. It is denoted by `Abs m` and defines a mapping that, given an object, in an invocation, selects the most specific method of the object having name `m` and the right types for the arguments of that invocation. This yields the solution, in our extended language, in Figure 2: `Map` has one *p_method* `F` whose type is `Fun → Object`. This means that the methods bound

```

public class FList {
    private Object elem;
    private FList next;
    public FList () {elem=null; next=null;}
    public FList Insert (Object x) {
        FList l= new FList(); l.elem=x; l.next=this; return l; }
    public Object Val () { return elem; }
    public FList Tail () { return next; } }
public abstract class Shape {
    public abstract Double Area();
    public abstract Double Perimeter();}
public class Circle extends Shape {
    private double radius;
    public Circle(double r){radius=r;}
    public Double Area() {return new Double(radius*radius*Math.PI);}
    public Double Perimeter() {return new Double(radius*2*Math.PI);} }
public class Rectangle extends Shape {
    private double base;
    private double height;
    public Rectangle(double b, double h){base=b; height=h;}
    public Double Area() {return new Double(base*height);}
    public Double Perimeter() {return new Double(2*(base+height));}}
public class FListShape extends FList {
    public FListShape Insert (Object x) {...}
    public FList FListArea(){
        FList L= new FList();
        if (Val()!=null) {L=((FListShape)Tail()).FListArea();
            L=L.Insert(((Shape) Val()).Area());}
        return L;}
    public FList FListPerimeter(){
        FList L= new FList();
        if (Val()!=null) {L=((FListShape) Tail()).FListPerimeter();
            L=L.Insert(((Shape) Val()).Perimeter());}
        return L; }}

```

Fig. 1. A class of geometric shapes in Java

to `F` have no arguments and compute a value of type `Object` as results. `FListShape` is still containing two methods, whose bodies are invocations of `Map: Map (Abs Area)` for `FListArea` and `Map (Abs Perimeter)` for `FListPerimeter`.

2.4 Integration with ordinary OO mechanisms

Since `p_method` is a function from objects (or classes) into instance (respectively, class) methods, in the evaluation of `L.Maps (Abs Area)`, for instance, the parameter `Abs Area` stands for the function that, given an object v , returns the *most specific* [GJSB00] method `Area` defined for v . Assumed that `L` is the list (v_1, \dots, v_n) , then `L.Map (Abs Area)` computes the list (r_1, \dots, r_n) where, for each $1 \leq i \leq n$, r_i is the result of $v_i.m_i()$ and m_i is the *most specific* method `Area` of object v_i . For $1 \leq i \neq j \leq n$, the methods m_i and m_j selected for object v_i and v_j , respectively, may differ, as it is the case when v_i is an instance of class `Circle` while v_j is instance of class `Rectangle`. Though objects of class `Circle` and objects of class `Rectangle` are objects of the superclass `Shape` and they may inherit from `Shape` the code for many methods, the code of method `Area` is based on subclass specialization and is different in the two classes. In this way we obtain a perfect integration of higher order generalization with the mechanisms of class hierarchy and inheritance allowing to write programs of an OO language using a higher order methodology .

3 Extension to the language

In this section we discuss the extensions to Java to support higher order generalization, in particular parameter passing and method invocation.

3.1 Method declaration

The syntax for method declaration, [GJSB00] §8.4, is modified in the following way:

$$\begin{aligned} \textit{MethodDeclaration}: & \quad \textit{ResultType Identifier} ([\textit{FType Identifier} \{, \textit{FType Identifier}\}]) \\ & \quad \textit{Block} \\ \textit{FType}: & \quad \textit{Type} \mid \textit{Fun FTLList} \rightarrow \textit{Type} \\ \textit{FTList} ::= & [\textit{FType} \{, \textit{FType}\}] \end{aligned}$$

it defines a new syntactic category, *FType*, which replaces *Type* in *MethodDeclaration*. *FType* can be an ordinary *Type* or a newly defined type, identified by the type constructor `Fun` that specifies the types of the arguments and type of the result of the methods that can be bound to the parameter, see §3.2.a.

3.2 Method invocation

The syntax for invocation, [GJSB00] §15.12, is modified in the following way:

$$\begin{aligned} \textit{Expression}: & \quad \textit{Expression.Identifier} ([\textit{AExpression} \{, \textit{AExpression}\}]) \\ & \quad \mid \textit{Expression.Parameter} ([\textit{AExpression} \{, \textit{AExpression}\}]) \\ \textit{AExpression}: & \quad \textit{Expression} \\ & \quad \textit{Abs Identifier} \end{aligned}$$

with the following meanings and constraints:

- (a) method invocation is extended by the second rule of *Expression*. It adds invocations of the form $e.p(l)$ where p is a formal p_method and e and l are object and list of arguments, respectively, of the invocation. Let $Abs\ m$ be the actual p_method bound to p . Let p be of type $Fun\ t_1, \dots, t_n \rightarrow t$. Let $\llbracket e \rrbracket$ be the value computed by e . Then the meaning of $e.p(l)$ is the invocation of $e.m(l)$, provided that *i*) l is a list of expressions of type t_1, \dots, t_n , and *ii*) m is the name of a method, of the hierarchy of $\llbracket e \rrbracket$, that applies to a list of arguments of type t_1, \dots, t_n and returns a value of type t .
- (b) actual parameters are extended by the second rule of *AExpression*. It adds parameters of the form $Abs\ m$, where m is the name of a method. The meaning is a function that given an object c of any class hierarchy C yields the most specific method of c having name m .

3.3 Example

In the extended language we can write, for instance, in a class C , the declaration:

```
int A (B X, Fun int → int F) {return X.F(0);}
```

This defines a method A , which returns an integer and has two parameters: X of type B and F of type $Fun\ int \rightarrow int$. Let $Abs\ M$ be the actual parameter bound to F . Then, in the invocation $X.F(0)$, in the body of A , the most specific method of the class hierarchy of (the object bound to) X , having name M , one argument of type int and returning a value of type int , is selected and applied, with argument 0 , to X . For instance, let c be an object of class C , and b be an object of class B , the invocation $c.A(b, Abs\ M)$ applies A to c with arguments b and the p_method $Abs\ M$. In the evaluation of the body of A , $X.F(0)$ yields the invocation of the most specific method of b which has name M and type $int \rightarrow int$. Because of Java overriding many methods may exist with name M and type $int \rightarrow int$ for the objects of class B . Hence different objects b may involve different methods in the evaluation of $X.F(0)$.

4 Implementation

4.1 Preprocessing: Structure of a syntactic translation

The implementation proposed in this paper uses a preprocessing technique which maps programs, in the extended syntax, back into programs of Java 1.4 [GJSB00]. A formal definition of the transformation is given by the mapping \mathcal{E} . It applies, separately, to each class of the source program producing a corresponding class of an equivalent program in Java 1.4. \mathcal{E} is a compositional transformation and this allows to express the translation through the collection of rules of Figures 4 and 5 that are descending on the class syntactic structure of the extended language. Because of space limitations, we restrict the presentation to the class structure of Figure 3. \mathcal{E} is indexed by the additional parameter ρ : It is an environment for the scope information; It contains the binding $\langle name, FType \rangle$ of each formal p_method, visible in the code, currently transformed by \mathcal{E} . At the basis of \mathcal{E} there is the transformation of the higher order introduced by the functions $Abs\ m$ into structures that can be handled at first order, in Java. All these functions:

```

public class FList {
    ...same as above for instance variables, constructor and methods val, etc.
    public FList Map(Fun → Object F){
        FList L=new FList();
        if (Val()!=null){L=Tail().Map(F);
            L=L.Insert(Val().F());}
        return L; }}
public class FListShape extends FList{
    public FListShape Insert (Object x){...}
    public FList FListArea() {return Map(Abs Area); }
    public FList FListPerimeter() {return Map(Abs Perimeter); }}

```

Fig. 2. Classes FList and FListShape in the extended language

```

ClassDeclaration ::= class Identifier [extends Type] [implements TypeList] {{ MemberDecl} }
MemberDecl ::= ;
                ModifiersOpt FieldDeclarator
                ModifiersOpt Identifier FParameters [throws QualifiedIdentifierList] Block
                ModifiersOpt Type Identifier FParameters [throws QualifiedIdentifierList] Block
                ModifiersOpt void Identifier FParameters [throws QualifiedIdentifierList] Block
                ModifiersOpt ClassOrInterfaceDeclaration
                [static] Block
FParameters ::= [FParameter {, FParameter}]
FParameter ::= [ final ] FType VariableDeclaratorId
FType ::=      Type |Fun FTList → Type                |Fun FTList → void
FTList ::= [FType {, FType}]
Selector ::=   .Identifier [Arguments]                 |.Par Arguments                 |.this
                |.super SuperSuffix                   |.new InnerCreator |[Expression]
Arguments ::= ([AExp {, AExp})
AExp ::= Expression | Abs Ide

```

Fig. 3. Extended syntax [GJSB00]

```

Let ClassDef ≡ public class A {
    ModifiersOpt Type0 Ide0 [=Exp0];
    ...
    ModifiersOpt Typeh Ideh [=Exph];
    ModifiersOpt A (FParameterC0)BlockC0
    ...
    ModifiersOpt A (FParameterCk)BlockCk
    ModifiersOpt TypeM0 IdeM0 (FParameterM0) BlockM0
    ...
    ModifiersOpt TypeMk IdeMk (FParameterMk) BlockMk
    ModifiersOpt void IdeMk+1 (FParameterMk+1) BlockMk+1
    ...
    ModifiersOpt void IdeMn (FParameterMn) BlockMn
 $\mathcal{E}[\text{ClassDef}]_\rho = \text{public class A implements ApplyClass} \{$ 
    ModifiersOpt Type0 Ide0 [=Exp0];
    ...
    ModifiersOpt Typeh Ideh [=Exph];
    ModifiersOpt A (FParameterC0)BlockC0
    ...
    ModifiersOpt A (FParameterCk)BlockCk
    ModifiersOpt TypeM0 IdeM0 ( $\mathcal{E}[FParameter_{M_0}]_\rho$ ) $\mathcal{E}[Block_{M_0}]_{\rho'_0}$ 
    ...
    ModifiersOpt TypeMk IdeMk ( $\mathcal{E}[FParameter_{M_k}]_\rho$ ) $\mathcal{E}[Block_{M_k}]_{\rho'_k}$ 
    ModifiersOpt void IdeMk+1 ( $\mathcal{E}[FParameter_{M_{k+1}}]_\rho$ ) $\mathcal{E}[Block_{M_{k+1}}]_{\rho'_{k+1}}$ 
    ...
    ModifiersOpt void IdeMn ( $\mathcal{E}[FParameter_{M_n}]_\rho$ ) $\mathcal{E}[Block_{M_n}]_{\rho'_n}$ 

```

Fig. 4. Transformation \mathcal{E} - part 1

- differ one another for the method that must be selected once the object to apply to and the types of the arguments of the invocation are known, while
- share the computation structure *i)* to apply the function to the object and the types, *ii)* to access the class hierarchy of the object, *iii)* to find the most specific method with that name and types of the arguments, *iv)* to apply the selected method to the object with the arguments of the invocation.

The idea is to have the computation structure of those functions as a sort of run time support that is included in the classes of the transformed programs, and used through suitable methods. This leads to methods, `Apply`, `ApplyS` that deal with all phases *i)-iv)* and are specific to each source class having methods that are passed as `p_method` in the source program. Invocations of `Apply` and `ApplyS` replace, in the transformed program, invocations in which the invoked method is a `p_method`. An interface `ApplyClass` defines them as two abstract methods:

```
public interface ApplyClass {
    public abstract Object Apply(String M, Object [] Pars);
    public abstract void ApplyS(String M, Object [] Pars);}
```

They have two parameters:

1. `M` is the string univocally (see *ToS* below) associated to the name of the method to invoke,
2. `Pars` is the array of the actual parameters of the method to invoke. Each parameter of the array is cast to the corresponding type of the argument type list bound to the `p_method` in the environment ρ .

`Apply` and `ApplyS` find, through `Dispatcher`, the method whose name is equal to `M` and invoke it with appropriate parameters taken from `Pars`. They differ because:

- `Apply` is used for methods returning a value, i.e. methods whose invocation is an expression. In this case `Apply` returns an object: the one computed by the invoked method. A type cast in \mathcal{E} is imposed on `Apply` result.
- `ApplyS` is used for methods which do not return any value. In this case `ApplyS` simply invokes the method. No type cast is required.

Auxiliary method `Dispatcher` is introduced to help in structuring phase *iii)*. It traverses the class hierarchy to find the method of a given name. Found the class containing the right definition, it computes the position in the class of the method. A string, namely `ToS(Idei)`, is univocally associated to each method name, `Idei`, for this purpose.

4.2 Example: The transformed program

The program resulting from the preprocessing of the program in Figure 2 is presented in Figures 6,7 and 8. The code for the class `Rectangle` is omitted since it is equal to the one for `Circle`. All environments ρ involved are empty except for class `FList`.

Let us consider now the execution of the method `FListArea` on a list of `Shapes` constituted of two elements a circle *c* and a rectangle *r*. Suppose the circle is the first element in the list, obtained applying `Val` in the first `Map` invocation. Hence, applying `Apply` to such element (of the class `Circle`) with argument the string "Area", yields the invocation of the method `Area` of the class `Circle`. In the second recursive invocation, a

rectangle is obtained by `Val` invocation. Hence applying `Apply` to such element (of the class `Rectangle`) with still `"Area"` yields the invocation of the method `Area` of the class `Rectangle`. So two different methods `Area` of the class `Circle` and `Area` of the class `Rectangle` are invoked, within a `Map` invocation with the same `p_method F`. This would not be possible if `F` value were a method, but more important, this is the behaviour required in a language with class hierarchy and inheritance mechanisms.

4.3 Syntactic simplifications

We confined the presentation of \mathcal{E} in several ways. However, \mathcal{E} can be completed for full Java 1.4 extended with the mechanisms for `p_methods`. In particular, source programs may contain: *i)* `p_methods` that yield methods with more than one argument or that are class methods, *ii)* exceptions, *iii)* class hierarchies involving abstract classes and interfaces, *iv)* inner, embedded and local classes.

4.4 Overloaded methods

Currently, \mathcal{E} assumes that overloaded methods are not yielded by `p_methods`. In fact, method overloading is solved, in Java, at static time by selecting, in each invocation, the more specific method among those matching the given method signature. In our approach the signature of the method yielded by a `p_method` is known at run time.

4.5 Program locality

Our approach maintains textual locality [DNR06]. In particular: *i)* the extended language does not modify the Java class structure; *ii)* the transformation \mathcal{E} does not modify the class structure of the source program. Noting that *ii)* cannot be obtained when *function pointers* [McC01] are used since they require the introduction, in the program, of suitable additional classes to contain the function pointer definitions.

4.6 Types and static type checking

The type of a formal `p_method` is `Fun FTList → Type`: It states number and types of the arguments and type of the result, of the methods that can be used in the invocation in which the `p_method` applies. It has the advantage that such methods belong to different, possibly unrelated, class hierarchies. The effective hierarchy, as well as the effective definition, for overridden methods, depends on the object the `p_method` applies to, in the invocation. This enhances flexibility but forbids the ability to check, at compile time, the involved hierarchies for the effective presence of a method with the right *signature* [GJSB00]. In this case, the transformed program can throw a `MethodNotFoundException()`.

5 Conclusions

The paper discussed a preprocessing implementation that supports Java 1.4 extended with higher order functionalities. The approach is based on `p_methods`, which introduce functions that map each object (class) of the program into the most specific method of

the object (resp. class) having a given name and type. This notion offers a restricted, disciplined, form of function abstraction but suitable to the integration of higher order and object oriented programming. The main advantages of this solution are that it is simple and expressive, it copes with code reusability and it is well suited to class hierarchy and inheritance mechanisms of Java. It is reasonably efficient, since it uses preprocessing. Eventually it does not require any modification to JVM.

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```

private static int Dispatcher(String S){int pos=-1;
    if (S.equals(ToS(IdeM0))) pos=0;
    else if ...
    else if (S.equals(ToS(IdeMk))) pos=k;
    else if (S.equals(ToS(IdeMk+1))) pos=k+1;
    else if ...
    else if (S.equals(ToS(IdeMn))) pos=n;
    return pos;}
public Object Apply(String M, Object Par) throws MethodNotFoundException{
    int pos=Dispatcher(M);
    switch (pos){case 0 : return IdeM0(Par);
        ...
        case k : return IdeMk(Par);}
    throw new MethodNotFoundException();}
public void ApplyS(String M, Object Par) throws MethodNotFoundException {
    int pos=Dispatcher(M);
    switch (pos){case k+1 : IdeMk+1(Par);
        ...
        case n : IdeMn(Par);}
    throw new MethodNotFoundException();}

```

where: $\rho'_i = \mathcal{R}[\text{FParameter}_{M_k}]_\rho$

$$\mathcal{E}[\text{FParameter}]_\rho = \begin{cases} \text{Type Ide} & \text{with FParameter} = \text{Type Ide} \\ \text{String Ide with FParameter} = \text{Fun FType} \rightarrow \text{Type Ide} \\ \text{String Ide with FParameter} = \text{Fun FType} \rightarrow \text{void Ide} \end{cases}$$

$\mathcal{R}[\text{FP}]_\rho(\text{ide}) = \langle \text{Fun FType Type} \rangle$ if $\text{FP} = \text{Fun FType} \rightarrow \text{Type Ide}$

$\mathcal{R}[\text{FP}]_\rho(\text{ide}) = \langle \text{Fun FType Type} \rangle$ if $\text{FP} = \text{Fun FType} \rightarrow \text{void Ide}$

$\mathcal{R}[\text{FP}]_\rho(x) = \rho(x)$ if $\text{FP} = \text{Type Ide}$

$\mathcal{E}[\text{Block}]_\rho = \mathcal{E}[\text{St}]_\rho; \mathcal{E}[\text{StList}]_\rho$ with $\text{Block} = \text{St}; \text{StList}$

$\mathcal{E}[\text{Arguments}]_\rho = [\mathcal{E}[\text{AExp}]_\rho\{\}, \mathcal{E}[\text{AExp}]_\rho\{\}]$

$\mathcal{E}[\text{AExp}]_\rho = \mathcal{E}[\text{Expression}]_\rho \mid \text{ToS}(\text{Ide})$

$$\mathcal{E}[\text{St}]_\rho = \begin{cases} ((\text{ApplyClass})\mathcal{E}[\text{Exp}_1]_\rho).\text{ApplyS}(\text{Par}, (\text{FType})\mathcal{E}[\text{Exp}_2]_\rho), & \text{with St} = \text{Exp}_1.\text{Par}(\text{Exp}_2) \wedge \\ \mathcal{E}[\text{Exp}_1]_\rho.\text{Ide}(\mathcal{E}[\text{Exp}_2]_\rho), & \rho(\text{Par}) = \langle \text{Fun FType} \rightarrow \text{void} \rangle \\ \text{if}(\mathcal{E}[\text{Exp}]_\rho)\mathcal{E}[\text{St}_1]_\rho \text{ else } \mathcal{E}[\text{St}_2]_\rho; & \text{with St} = \text{Exp}_1.\text{Ide}(\text{Exp}_2) \wedge \\ \text{while}(\mathcal{E}[\text{Exp}]_\rho)\mathcal{E}[\text{St}]_\rho & \rho(\text{Ide}) = \perp \\ \text{etc.} & \text{with St} = \text{if Exp St}_1 \text{ else St}_2 \\ & \text{with St} = \text{while Exp St} \end{cases}$$

$$\mathcal{E}[\text{Exp}]_\rho = \begin{cases} (\text{Type})((\text{ApplyClass})\mathcal{E}[\text{Exp}]_\rho).\text{Apply}(\text{Par}, (\text{FType})\mathcal{E}[\text{Exp}]_\rho), & \text{with Exp} = \text{Exp}.\text{Par}(\text{Exp}) \wedge \\ \mathcal{E}[\text{Exp}]_\rho.\text{Ide}(\mathcal{E}[\text{Exp}]_\rho), & \rho(\text{Par}) = \langle \text{Fun FType} \rightarrow \text{Type} \rangle \\ \mathcal{E}[\text{Exp}]_\rho \text{ Op } \mathcal{E}[\text{Exp}]_\rho & \text{with Exp} = \text{Exp}.\text{Ide}(\text{Exp}) \wedge \\ \text{etc.} & \rho(\text{Ide}) = \perp \\ & \text{with Exp} = \text{Exp Op Exp} \end{cases}$$

Where:

- Exp , St , StList , FParameters , Ide stand for *Expression*, *Statement*, *StatementList*, *Formal-Parameters*, *Identifier* respectively;
- $\text{ToS}(m)$ computes the string univocally associated to the method of name m

Fig. 5. Transformation \mathcal{E} - part 2

```

public class FList implements ApplyClass {
    ...same as Figure 1 for instance variables etc.
    public FList Map(String F){
        FList L=new FList();
        if (Val()!=null) { L=Tail().Map(F);
            L=L.Insert((Object)((ApplyClass)Val().Apply(F,empty))); }
        return L; }
    private static int Dispatcher(String S){
        int pos;
        if (S.equals("Insert")) pos=0;
        else if (S.equals("Val")) pos=1;
        else if (S.equals("Tail")) pos=2;
        else if (S.equals("Map")) pos=3;
        else pos=-1;
        return pos;}
    public Object Apply(String M, Object [] Pars) throws MethodNotFoundException{
        int pos=Dispatcher(M);
        switch (pos){ case 0:return Insert(Pars[0]);
                    case 1:return Val();
                    case 2:return Tail();
                    case 3:return Map(Pars[0]);}
        return new MethodNotFoundException();}
    public void ApplyS ...

```

where: $\rho = \{ \langle F, \text{Fun} \rightarrow \text{Object} \rangle \}$; *empty* is a vector variable containing the parameters of the method `Object [] empty={}`

Fig. 6. Class FList

```

public class Circle extends Shape implements ApplyClass {
    ...same as Figure 1 for instance variables etc.
    private static int Dispatcher(String S){int pos;
        if (S.equals("Area")) pos=0;
        else if (S.equals("Perimeter")) pos=1;
        else pos=-1;
        return pos;}
    public Object Apply(String M, Object [] Pars) throws MethodNotFoundException{
        int pos=Dispatcher(M);
        switch (pos){ case 0:return Area();
                    case 1: return Perimeter();}
        return new MethodNotFoundException();}
    public void ApplyS ...

```

Fig. 7. Class Circle

```

public class FListShape extends FList implements ApplyClass {
    public FList FListArea{ return Map("Area");}
    public FList FListPerimeter{return Map("Perimeter");}
    private static int Dispatcher(String S){
        int pos;
        if (S.equals("FListArea")) pos=0;
        else if (S.equals("FListPerimeter")) pos=1;
        else pos=-1;
        return pos;}
    public Object Apply(String M, Object [] Pars){
        int pos=Dispatcher(M);
        switch (pos){ case 0: return FListArea();
                    case 1: return FListPerimeter();}
        return new MethodNotFoundException();}
    public Object ApplyS ...

```

Fig. 8. FListShape