

Let us take  $n \ l = \text{match } (n, l) \text{ with}$

$$\left\{ \begin{array}{l} (\emptyset, l) \rightarrow [] \\ (n, [ ]) \text{ when } \underline{n > 0} \rightarrow [] \\ (n, x :: xs) \text{ when } n > 0 \rightarrow x :: (\text{take } (n-1) \ xs) \end{array} \right.$$

Let us drop  $n \ l = \text{match } (n, l) \text{ with}$

$$\left\{ \begin{array}{l} (\emptyset, l) \rightarrow l \\ (n, [ ]) \text{ when } \underline{n > 0} \rightarrow [] \\ (n, x :: xs) \text{ when } n > 0 \rightarrow \text{drop } (\underline{n-1}) \ \underline{xs} \end{array} \right.$$

$\forall n \in \mathbb{N}, l \in \text{'a list}. (\text{take } n \text{ } l) @ (\text{drop } n \text{ } l) = l$

$\mathbb{N} * \text{'a list} : \text{domains}$

$$(n, l) \sqsubseteq (n', l') \equiv$$

$$\begin{cases} n' \neq \emptyset \wedge n = n' - 1 \\ l' \neq [] \wedge l = \text{tl } l' \end{cases} \wedge$$

CASO BASE 1

(take  $\alpha(l)$ ) @ (drop  $\alpha(l)$ )

= { def. di take e drop, 1° patt. }

[] @ l

= { def. di @ }

l

CASO BASE 2

(take  $n[j]$ ) @ (drop  $n[j]$ )

= { ipotesi:  $n \in \mathbb{N}$ ,  $n > 0$ , 2° patt. }

[] @ []

= { def. di @ }

[]

CASO INDUTTIVO

Ip:  $m > 0$

$$\boxed{(\text{take } (m-1) \text{ xs}) @ (\text{drop } (m-1) \text{ xs}) = \text{xs}}$$

IP. INDUTTIVA

$$\Rightarrow (\text{take } n \text{ x::xs}) @ (\text{drop } n \text{ x::xs}) = x :: \text{xs}$$

$$(\text{take } n \text{ x::xs}) @ (\text{drop } n \text{ x::xs})$$

$$= \{ \text{Ip. } m > 0, 3^{\circ} \text{ pattern di take e drop} \}$$

$$(x :: (\text{take } (m-1) \text{ xs})) @ (\text{drop } (m-1) \text{ xs})$$

$$= \{ \text{proprietà } (x :: \text{xs}) @ \text{ys} = x :: (\text{xs} @ \text{ys}) \}$$

$$x :: \underline{((\text{take } (m-1) \text{ xs}) @ (\text{drop } (m-1) \text{ xs}))}$$

$$= \{ \text{Ip. induuttiva } (m-1, \text{xs}) \subset (n, \text{xs}) \}$$

$$x :: \text{xs}$$

# REVERSE (rovesciare) di una lista

1)  $\text{rev} [] = []$

~~...  
...  
x~~

2)  $\text{rev} [x_1; x_2; \dots; x_n] = \underbrace{\text{rev} [x_2; \dots; x_n]}_{[x_n; x_{n-1}; \dots; x_2]} @ [x_1]$

let rec rev l = match l with

$$[] \rightarrow []$$

$$| x :: xs \rightarrow (\text{rev } xs) @ [x];;$$

$\text{rev} : 'a \text{ list} \rightarrow 'a \text{ list}$

$\text{rev } [1; 2; 3]$

=  $\{ 2^{\circ} \text{ pattern di rev} \}$

$\underline{(\text{rev } [2; 3])e [1]}$

=  $\{ 2^{\circ} \text{ pattern di rev} \}$

$((\text{rev } [2])e [2])e [1]$

=  $\{ 2^{\circ} \text{ pattern di rev} \}$

$((\text{rev } [])e [3])e [2]e [1]$

=  $\{ 1^{\circ} \text{ pattern di rev} \}$

$([]e [3])e [2]e [1]$

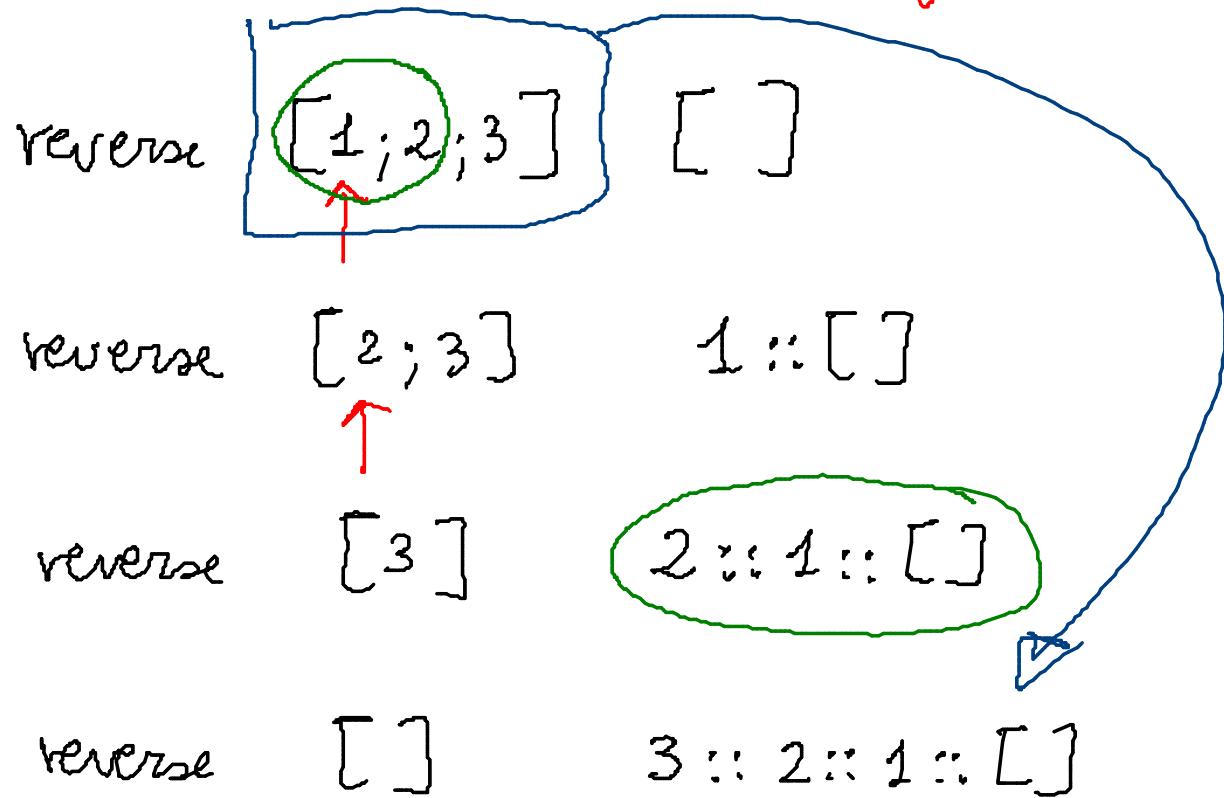
=  $\{ \text{calcoli} \}$

$[3; 2; 1]$

## REVERSE CON ACCUMULATORE

reverse l **l1** accumulo posso passo il risultato  
che voglio ottenere

quando l si scrive , il secondo  
argomento rappresenta il risultato del calcolo



let rec reverse  $l$   $l1 =$  match  $l$  with  
 $[] \rightarrow l1$   
 |  $x::xs$   $\rightarrow$  reverse  $xs$   $x::l1$ ;;

reverse : 'a list  $\rightarrow$  'a list  $\rightarrow$  'a list

Relazione fra precedente insieme

$$(xs, x\underline{::} l1) \sqsubset (x::xs, \underline{l1})$$

$$\forall l, l'. (\text{reverse } l \ l') = (\text{rev } l) \circ l'$$

# let rev1  $l =$  reverse  $l$  [];;

rev1 : 'a list  $\rightarrow$  'a list

# FUNZIONI DI ORDINE SUPERIORE

Quantificazione "universale"  $\forall x \in l. p(x)$

let rec forall  $p$  l = match l with

[ ]  $\rightarrow$  true

|  $x :: xs \rightarrow (p x) \& (\text{forall } p xs)$  ; ;

forall :  $(\underline{a} \rightarrow \text{bool}) \rightarrow \underline{\underline{a}} \text{ list} \rightarrow \text{bool}$

$p$  è un predicato, cioè  
una funzione che restituisce un booleano

# let  $\text{mago} x = x > 0$  ;;

$\text{mago} : \text{int} \rightarrow \text{bool} = \langle \text{fun} \rangle$

# forall  $\text{mago} [1; 3] ;;$

- :  $\text{bool} = \text{true}$

# forall  $\text{mago} [-1; 1; \dots; 1000000]$

- :  $\text{bool} = \text{false}$

Calcolo  $(\text{mago} -1) \& (\text{mago} 1) \& \dots \& (\text{mago} 1000000) \& \text{true}$

= **false**

$$\begin{aligned} &= \text{forall } \text{mago} [1; 3] \\ &= \{ \text{2}^0 \text{ path} \} \\ &= (\text{mago} 1) \& (\text{forall } p [3]) \\ &= \{ \text{2}^0 \text{ path.} \} \\ &= (\text{mago} 1) \& (\text{mago} 3) \& \text{forall } \text{mago} [] \\ &= \{ \text{1}^0 \text{ path.} \} \\ &= (\text{mago} 1) \& (\text{mago} 3) \& \text{true} \\ &= \{ \text{def. di mago} \} \\ &= \text{true} \end{aligned}$$

let rec forall p l = match l with  
[] → true

|  $x :: xs$  when not( $p x$ ) → false

|  $x :: xs$  when ( $p x$ ) → forall p  $xs;;$

○ ————— ○

QUANTIFICATORE ESISTENZIALE

$\exists x \in l. p(x)$

let rec exists p l = match l with

[] → false

|  $x :: xs \rightarrow (p x) \text{ or } (\exists \text{exists } p xs);;$

exists: ('a → bool) → 'a list → bool = <fun>

Let rec exists  $\lambda l =$  match  $l$  with

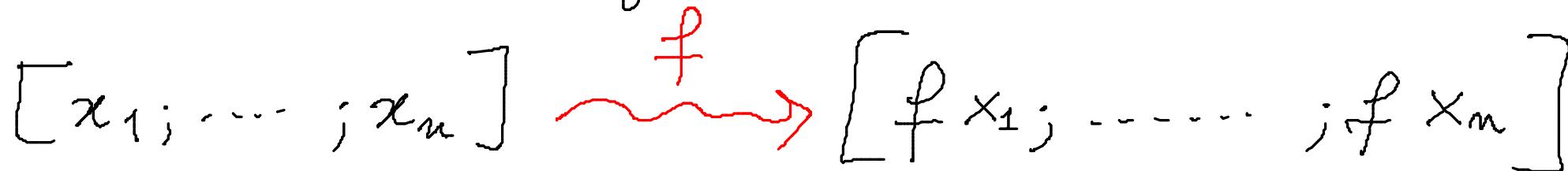
$[] \rightarrow \text{false}$

$x :: xs$  when  $(\lambda x) \rightarrow \text{true}$

$x :: xs$  when  $\text{not } (\lambda x) \rightarrow \text{exists } \lambda xs ;$



MAP : applicare una stessa operazione  
a tutti gli elementi di una sequenza



let rec map  $f$  l = match l with

$(f, x) ::$   
 $(f, x :: xs)$

$$\begin{cases} [] & \rightarrow [] \\ x :: xs & \rightarrow (fx) :: (\text{map } f xs) \end{cases}$$

map :  $('a \rightarrow 'b) \rightarrow 'a \text{ list} \rightarrow 'b \text{ list}$   
 $f$   $l$   $'b \text{ list}$   
isra.

Funzione che resterà tutti gli elementi negativi di una lista di interi

$$\text{atteraneg } [-1; 2; 3; -5; 4] = [\emptyset; 2; 3; \emptyset; 4]$$

let rec aterraneq l = match l with

$$[ ] \longrightarrow [ ]$$

$$\left| \begin{array}{lll} x :: xs & \text{when } x >= \emptyset & \rightarrow x :: (\text{atteraneq } xs) \end{array} \right.$$

$$\left| \begin{array}{lll} x :: xs & \text{when } x < \emptyset & \rightarrow \emptyset :: (\text{atteraneq } xs) \end{array} \right.$$

o ————— o

atteraneq: int list  $\rightarrow$  int list

# let g x = if x < 0 then 0 else x ;;

g : int → int

# let average l = map g l;;

average :  $\frac{\text{list}}{\text{int}} \rightarrow \frac{\text{list}}{\text{int}}$

map :  $(\frac{a}{\text{list}} \rightarrow \frac{b}{\text{list}}) \rightarrow [\frac{a}{\text{list}} \rightarrow \frac{b}{\text{list}}]$

g : int → int

DEFINIZIONI LOCALI    let ... in

# let ottomaneg l = let g x = if x < 0 then 0 else x  
              in  
              map g l ;;

# g 3 ;;  
^  
unbound

# ottomanag [-1 ; 2 ; -3] ;;  
- : int list = [0 ; 2 ; 0]

let E in E'



i nomi definiti in  
E sono VISIBILI in  
E' ma non altrove

# let y = let x = 10 in x + 1;

y : int = 11

# x + 2 ;  
  ^

unbound

# y + 2 ;  
  ^ ;  
- int = 13

FILTER : cancella da una sequenza  
tutti gli elementi che non soddisfano  
una data proprietà

let rec filter p l = match l with

$$[ ] \rightarrow [ ]$$

|  $x :: xs$  when (p x)  $\rightarrow$   $x :: (\text{filter } p xs)$

|  $x :: xs$  when not (p x)  $\rightarrow$   $(\text{filter } p xs) :: xs$

filter :  $(\alpha \rightarrow \text{bool}) \rightarrow 'a \text{ list} \rightarrow 'a \text{ list}$

# filter mago [ -1; 2; 3; -5 ] ;;

- : int list = [ 2; 3 ]

# let mago x =  $x > 0$  ;;