

# Models of Computation

Midterm Exam on April 9, 2010

## Exercise 1 (12)

Extend IMP with: (i) the new syntactic category  $S$  (with variables  $s$ ) of *expressions with side effects*, shortly *elat*, having  $\langle s, \sigma \rangle \rightarrow (\sigma', n)$  as well formed formulas for the operational semantics and semantic domain  $\Sigma \rightarrow (\Sigma \times N)_\perp$ ; and (ii) the new productions

$$s ::= c ; a \qquad c ::= x := s$$

with the following informal semantics: for the first production, the meaning is to return the memory produced by the command  $c$  as *modified memory* by elat  $s$ ; and the evaluation of the arithmetic expression  $a$  in such a memory as *result* of  $s$ . For the second production, the meaning is first to evaluate  $s$  and then to execute the assignment, employing the memory modified by  $s$  as memory, and the result of  $s$  as value to assign to  $x$ .

The denotational semantics is as follows:

$$\begin{aligned} \mathcal{S}[c; a]\sigma &= (\lambda\sigma'.(\sigma', \mathcal{A}[a]\sigma'))^*(\mathcal{C}[c]\sigma) \\ \mathcal{C}[x := s]\sigma &= \mathbf{case} \mathcal{S}[s]\sigma \mathbf{ of} : \\ &\quad \perp_{(\Sigma \times N)_\perp} : \perp_{\Sigma_\perp} \\ &\quad (\sigma', n) : \sigma'[n/x]. \end{aligned}$$

Define the operational semantics of the two new constructs and evaluate the command  $x := (x := 3; 5)$  according to both the operational and the denotational semantics. Prove the operational - denotational equivalence and show (employing either the operational or the denotational semantics) that  $x := (c ; a)$  is equivalent to  $c ; x := a$ .

## Exercise 2 (8)

Consider the context free (more precisely, right linear) grammar  $S ::= \alpha S \mid \beta$ , where the terminal alphabet is  $V = \{a, b\}$ , with  $\alpha, \beta \in V^*$ . Derive the inference rules  $R$  corresponding to such a grammar and prove (in both ways) that the set of theorems is  $I_R = \{\alpha^n \beta \in S \mid n \in \omega\}$ .

Then consider the immediate consequences operator  $\hat{R}$  and check that  $I_R$  is its minimal fixpoint. Finally, discuss if other fixpoints exist for particular values of  $\alpha$  and  $\beta$ .

## Exercise 3 (10)

The *binomial coefficients*  $\binom{n}{k}$ , with  $n, k \in \omega$  e  $0 \leq k \leq n$ , are defined by:

$$\binom{n}{0} = \binom{n}{n} = 1 \qquad \binom{n+1}{k+1} = \binom{n}{k} + \binom{n}{k+1}.$$

Prove that the above definition is given by well-founded recursion, specifying the well-founded relation and the corresponding function  $F(b, h)$ .

Then consider the HOFL program:

$$t = \mathit{rec} f. \lambda n. \lambda k. \mathbf{if} \ k \ \mathbf{then} \ 1 \ \mathbf{else} \ \mathbf{if} \ n - k \ \mathbf{then} \ 1 \ \mathbf{else} \ ((f \ n - 1) \ k) + ((f \ n - 1) \ k - 1),$$

compute its type and evaluate the canonical form of the term  $((t \ 2) \ 1)$ .